

# **Notas Breves**



ON FUNCTIONS THAT CANNOT BE MV-TRUTH VALUES  
IN ALGEBRAIC STRUCTURES

ENRIC TRILLAS

ABSTRACT

*It is shown, in a general frame and playing with idempotency, that in order to have on a given lattice a Multiple Valued Logic preserving the lattice structure, the only t-norms and t-conorms allowing to modelize the truth values of  $a \vee b$ ,  $a \wedge b$  and  $a \rightarrow b$  are Min and Max, respectively, apart from ordinal sums.*

Let  $(L, *)$  be an algebraic structure, that is a non-empty set  $L$  with a binary operation  $L \times L \rightarrow L, *(a, b) = a * b$ , and let  $V$  be a function  $L \rightarrow [0, 1]$  "evaluating" the elements of  $L$  in some predetermined sense (a certainty factor in the language of Expert Systems) such that, for all  $(a, b) \in L \times L$ ,

$$V(a * b) = T(V(a), V(b)) \tag{1}$$

being  $T$  a t-norm [1].

**Theorem 1.** Any function  $V : L \rightarrow [0, 1]$ , verifying (1), preserves the relation  $a \leq b$  given in  $L$  by  $a * b = a$ .

*Proof.* If  $a \leq b$ , as  $T \leq \text{Min}$  [1], we have  $V(a) = V(a * b) \leq \text{Min}(V(a), V(b)) \leq V(b)$ .

**Theorem 2.** If there exists  $z \in L$  such that  $z = z * z$  and  $V(z) \in (0, 1)$ , then the only t-norms  $T$  for which (1) is possible are Min and the ordinal sums.

*Proof.* As  $V(z) = V(z * z) = T(V(z), V(z))$ , we see that  $V(z)$  is an idempotent for  $T$ , so that  $T$  has to be [1] either Min or an ordinal sum.

Theorems 1 and 2 can be used when  $L$  is an lower-semilattice [2] with  $* = \wedge$ , the meet operation. In that case, functions  $V : L \rightarrow [0, 1]$  verifying  $V(a \wedge b) = T(V(a), V(b))$  and effectively ranging in  $(0, 1)$  are candidates to be Multiple-Valued Truth Values, but with the strong limitation of  $T$  being only Min or an ordinal Sum. The same happens if  $L$  is a set of propositions equipped with a sort of conjunction and with a certainty factor  $V$  [3], when there exists just one proposition  $p$  such that  $p \wedge p = p$  and  $0 < V(p) < 1$ .

Similar proofs are valid for the next theorems.

**Theorem 3.** If there exists  $z \in L$  such that  $z * z = z$  and  $V(z) \in (0, 1)$ , then the only t-conorms  $S$  for which it holds

$$V(a * b) = S(V(a), V(b)) \quad (2)$$

for all  $(a, b) \in L \times L$  are the Max and the ordinal sums.

**Theorem 4.** Any function  $V : L \rightarrow [0, 1]$ , verifying (2) preserves the relation  $a \leq b$  given in  $L$  by  $a * b = b$ .

Theorems 3 and 4 can be used when  $L$  is an upper-semilattice [2] with  $* = \vee$  the join operation. In such case, functions  $V$  verifying  $V(a \vee b) = S(V(a), V(b))$  and effectively ranging in  $(0, 1)$  are also candidates to be MV-Truth Values, but with the limitation of  $S$  being either Max or an ordinal sum. The same happens if  $L$  is a set of propositions equipped with some disjunction  $\vee$  and a certainty factor  $V$  [3], when there exists just one proposition  $p$  such that  $p \vee p = p$  and  $0 < V(p) < 1$ .

Consequently, for lattices  $(L, \vee, \wedge)$ , the only possibilities for representing actual MV-Truth Values using t-norms or t-conorms by means of (1) or (2) respectively, are  $T = \text{Min}$  or  $S = \text{Max}$  and, in both cases, ordinal sums. Other t-norms or t-conorms will give only

$V(a) \in \{0, 1\}$  for each  $a \in L$ , and no actual Multiple-Valued Logic will be available with such functions  $V$ . This result completes what is proven in [4], [5] and [6].

In the last case, and provided that  $\neg\neg a = a$  for each  $a \in L$ , it should be pointed out that it is excluded the possibility of giving the truth value of  $a \rightarrow b = \neg a \vee b$  by means of  $V(a \rightarrow b) = 1 - T(1 - V(\neg a), 1 - V(b))$ , where  $V(\neg a) = 1 - V(a)$ , by virtue of  $V(\neg a \rightarrow a) = V(a \vee a) = V(a) = 1 - T(V(\neg a), 1 - V(a)) = 1 - T(1 - V(a), 1 - V(a))$ , or  $1 - V(a) = T(1 - V(a), 1 - V(a))$ ; and, if  $V(a) \in (0, 1)$ ,  $T$  should be either Min or an ordinal sum.

Thus, it is clear that for compatibility with a lattice structure on  $L$  only  $T = \text{Min}$  or  $T = \text{an ordinal sum}$ , are admissible t-norms. Of course, it is possible to use other t-norms to modelize  $V(a \rightarrow b)$  only, as is the case of Lukasiewicz's Logic [7] in which  $V(a \vee b) = \text{Max}(V(a), V(b))$ ,  $V(a \wedge b) = \text{Min}(V(a), V(b))$  and  $V(a \rightarrow b) = \text{Min}(1, 1 - V(a) + V(b)) = 1 - L(V(a), 1 - V(b))$ , where  $L(x, y) = \text{Max}(0, x + y - 1)$  is the so-called Lukasiewicz t-norm [1]. This is also the case of Zadeh's Fuzzy Logic [8] with  $L = [0, 1]^{[0, 1]}$  the set of all fuzzy sets over  $[0, 1]$ , and where  $(\mu_A \cap \mu_B)(x) = \text{Min}(\mu_A(x), \mu_B(x))$ ,  $(\mu_A \cup \mu_B)(x) = \text{Max}(\mu_A(x), \mu_B(x))$  and  $(\mu_A \rightarrow \mu_B)(x) = \text{Min}(1, 1 - \mu_A(x) + \mu_B(x))$ , for each  $x \in [0, 1]$ ,  $\mu_A$  and  $\mu_B$  being fuzzy sets over  $[0, 1]$ . In those cases no idempotency "pathology" arises.

**Remark.**

(I) If  $a \rightarrow b = \neg a \vee (a \wedge b)$  and  $\neg a \vee a = 1$ , for each  $a \in L$  (Excluded Middle Principle) then, setting  $V(a \rightarrow b) = V(\neg a \vee (a \wedge b)) = 1 - T(V(a), 1 - V(b))$ , we have  $V(\neg a \vee a) = 1 = V(a \rightarrow a) = 1 - T(V(a), 1 - V(a))$  and  $T(V(a), 1 - V(a)) = 0$  for all  $a \in L$ . This excludes t-norms "like"  $T = \text{Min}$  or  $T = \text{Prod}$  if we want to preserve both the Excluded Middle Principle and the actual MV-Truth Value. For t-norms "like L" (see [9]),  $T = \phi^{-1} \circ L \circ (\phi \times \phi)$  where  $\phi$  is an automorphism of  $[0, 1]$ , we will have  $\phi^{-1}(L(\phi(V(a)), \phi(1 - V(a))) = 0$ , that

is  $L(\phi(V(a)), \phi(1 - V(a))) = 0$ , or  $\text{Max}(0, \phi(V(a)) + \phi(1 - V(a)) - 1) = 0$ . Thus,  $\phi$  must verify  $\phi(V(a)) + \phi(1 - V(a)) \leq 1$  for each  $a \in L$  (as is actually in the Lukasiewicz case for  $\phi = j$ , the identity function, so that both the Excluded Middle Principle and the MV-Truth Value [4] hold.

(II) Of course, it is always possible to redefine  $V(a \vee b)$  in lattices with least element 0, and for "measure" purposes, by setting  $V(a \vee b) = S(V(a), V(b))$  if  $a \wedge b = 0$ . In the case of orthomodular lattices [2] in which  $a \rightarrow b = a' \vee (a \wedge b)$ ,  $a'$  being the orthocomplement of  $a$  (for Boolean algebras we have  $a \rightarrow b = a' \vee b$ ) it is possible to write  $V(a \rightarrow b) = S(V(a'), V(a \wedge b))$  because  $a' \wedge (a \wedge b) = 0$ . Nevertheless, in all those cases we do not obtain what is usually known as a truth-value function or truth-table for logical connectives.

#### References.

- [1] Schweizer, B. and Sklar, A. (1983) *Probabilistic Metric Spaces*. North-Holland.
- [2] Birkhoff, G. (1979) *Lattice Theory*. AMS Colloquium Pubs. Third. Ed.
- [3] Smets, Ph., Mamdani, E.H., Dubois, D. and Prade, H. Ed. (1988) *Non-Standard Logics for Automated Reasoning*. Academic Press.
- [4] Trillas, E. (1988) "Tercer exclòs, funcionalitat i bivaluació", in *Actes del VIIè Congrés Català de Lògica*. 81-83.
- [5] Trillas, E., Alsina, C. and Valverde, L. (1982) "Do we need Max, Min and 1-j in Fuzzy Set Theory?" in *Fuzzy Set and Possibility Theory*, R.R. Yager, Ed., Pergamon Press, pp. 275-297.
- [6] Alsina, C., Trillas, E. and Valverde, L. (1983) "On some logical connectives for Fuzzy Sets Theory", in *Jour. Math. Ann. and its Appl.* pp. 15-26.
- [7] Ackerman, R.
- [8] Zadeh, L.A. (1984) "A Theory of Commonsense Knowledge", in *Aspects of Vagueness*. H.H. Skala, S. Termini and E. Trillas (ed.), Reidel, pp. 257-295.

- [9] Ovchinnikov, S. and Roubens, M. (1988) "On Strict Preference Relations" (preprint).

Facultad de Informática

Universidad Politécnica de Madrid.

