THE UNIQUENESS PROBLEM IN THE THEORY OF NUMERICAL DIVERGENT SERIES AND FORMAL LAWS OF CALCULUS II (*)

by

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§ 2. Identity of an algorithm with those of EULER and BOREL.

1. Definition of convergence and addition algorithms through the formal laws of calculus. Let

be the set of s_n numerical convergent sebe the set of the numerical convergent series quences $\sum_{n=0}^{\infty}a_n<\infty;$ $\lim s_n < \infty$; and λο the isomorphism existing between and the complex field $C=R^2$ with respect to the formal laws of calculs used in the elemental theory of convergent sequences convergent series The problem of the generalized limits 1 addition of divergent series should be attempted in our meaning in this other way [21, no, 1, 2 1, Chap, V, 339]. Problem A. To extend the isomorphism by means of another isomorphism between a set $s_\lambda \mathrel{\sqsupset} s_0 \; (\; s_\lambda \neq s_0)$ $S_{\sigma} \supset S_0 \ (S_{\sigma} \neq S_0)$

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and the same complex plan $\mathcal C$ with respect to certain formal laws, for instance, the laws a', b', c', d' or d', e' and f' a, b, c, d, e and f stated below in the weak form that will be used. Def. 1. We will call permanent method of addition convergence or more briefiy, algorithm of addition convergence any solution λ σ to problem A with respect to an system of formal laws. The algorithm will be distributive when, among these laws, there exists the a' and b' Def. 2. The series $\sum_{n=0}^{n=\infty} a_n \in S_{\sigma}$ sequences $\{s_n\} \in S_{\lambda}$ are called σ -summable; λ-convergent; and the number s homologous in the isomorphism is called σ-sum λ-limit of the serie $\sum_{n=0}^{\infty} a_n$ sequence $\{s_n\}$ and is written $\sigma \sum_{n=0}^{\infty} a_n = s$ λ -lim $s_n = l$ Def. 3. We say that an algorithm s (strictly) weaker than another λ or that λ' is (strictly) stronger than σ λ

sometimes written

$$\lambda \subset \lambda' \ (\lambda \subset \lambda', \ \lambda \neq \lambda') \qquad \qquad | \qquad \qquad \sigma \subset \sigma' \quad (\sigma \subset \sigma', \ \sigma \neq \sigma')$$

when it is

$$s_{\lambda} \subset s_{\lambda'} \ (s_{\lambda} \subset s_{\lambda'}, \ s_{\lambda} \neq s_{\lambda'}) \qquad \qquad | \qquad \qquad S_{\sigma} \subset S_{\sigma'} \ (S_{\sigma} \subset S_{\sigma'}, \ S_{\sigma} \neq S_{\sigma'})$$

Prop. 1. (Transitive property). If

$$\lambda \subset \lambda'$$
 and $\lambda' \subset \lambda''$, $\lambda \subset \lambda''$ $\sigma \subset \sigma'$ and $\sigma' \subset \sigma''$, $\sigma \subset \sigma''$

Proof. Evidently, if

$$s_{\lambda} \subset s_{\lambda'}$$
 and $s_{\lambda'} \subset s_{\lambda''}$, $s_{\lambda} \subset s_{\lambda''}$ | $s_{\sigma} \subset s_{\sigma'}$ and $s_{\sigma'} \subset s_{\sigma''}$, $s_{\sigma} \subset s_{\sigma''}$

Def. 4. If

$$\{\lambda_{\alpha}\}$$
 | $\{\sigma_{\alpha}\}$

is an algorithms family of

with a variable parameter a within a set F, we will say that the set or union of these algorithms

$$U (\lambda_{\alpha}, \alpha \in F)$$

$$U (\sigma_{\alpha}, \alpha \in F)$$

is (strictly) weaker than an algorithm

or that the latter is (strictly) stronger than the said set of algorithms, and may write $\mathbf{U}(\lambda_{\alpha}, \alpha \in F) \subset \lambda$, $(\mathbf{U}(\lambda_{\alpha}, \alpha \in F) \subset \lambda, \mathbf{U}(\lambda_{\alpha}, \alpha \in F) \neq \lambda) + \mathbf{U}(\sigma_{\alpha}, \alpha \in F) \subset \sigma$, $(\mathbf{U}(\sigma_{\alpha}, \alpha \in F) \subset \sigma, \mathbf{U}(\sigma_{\alpha}, \alpha \in F) \neq \sigma)$ when

$$\begin{array}{c|c} & \textbf{U} \; (\textbf{S}_{\pmb{\lambda}\pmb{\alpha}} \; \alpha \; \epsilon \; F) \; \textbf{C} \; \textbf{S}_{\pmb{\lambda}}, \\ (\textbf{U} \; (\textbf{S}_{\pmb{\lambda}\pmb{\alpha}}, \alpha \; \epsilon \; F) \; \textbf{C} \; \textbf{S}_{\pmb{\lambda}\pmb{\alpha}}, \; \textbf{U} \; (\textbf{S}_{\pmb{\lambda}\pmb{\alpha}}, \alpha \; \epsilon \; F) \; \neq \; \textbf{S}_{\pmb{\lambda}}) \end{array} \qquad \begin{array}{c|c} & \textbf{U} \; (\textbf{S}_{\pmb{\sigma}\pmb{\alpha}}, \alpha \; \epsilon \; F) \; \textbf{C} \; \textbf{S}_{\pmb{\sigma}}, \\ (\textbf{U} \; (\textbf{S}_{\pmb{\sigma}\pmb{\alpha}}, \alpha \; \epsilon \; F) \; \textbf{C} \; \textbf{S}_{\pmb{\alpha}}, \; \textbf{U} \; (\textbf{S}_{\pmb{\sigma}\pmb{\alpha}}, \alpha \; \epsilon \; F) \; \neq \; \textbf{S}_{\pmb{\sigma}} \end{array}$$

Def. 5. We will say that an

addition algorithm
$$\sigma$$
 | convergence algorithm λ

is equivalent to

a convergence algorithm,
$$\lambda$$
 $\;\;\;|\;\;\;\;$ an addition algorithm, σ

and we will write sometimes

when the two following conditions are verified

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma} \xrightarrow{} \{s_n\} \in S_{\lambda} \quad (1)$$

$$\begin{cases} s_n\} \in S_{\lambda} & \text{if } s_n = \infty \\ \sum_{n=0}^{\infty} a_n \in S_{\sigma} & \text{if } s_n = \infty \end{cases}$$

$$\lambda \lim s_n = \sigma \sum_{n=0}^{\infty} a_n \quad (2')$$

being

$$s_n = \sum_{n=0}^{n} a_n$$
 for $n = 0, 1, 2, ...$ (3) $a_0 = s_0$ and $a_n = \Delta s_{n-1}$ for $n = 1, 2, ...$ (3')

Prop. 2. (Converse property). If:
$$\sigma = \lambda \qquad | \qquad \lambda = \sigma$$
it is
$$\lambda = \sigma \qquad | \qquad \sigma = \lambda$$
Proof. If
$$(1) \text{ and } (2) \qquad | \qquad (1') \text{ and } (2')$$
are verified for the
$$\{s_n\} \qquad | \qquad (a_n)$$
defined in
$$(3), \qquad | \qquad (3'), \qquad (3'), \qquad (1) \text{ and } (2)$$
for the
$$\sum_{n=0}^{\infty} a_n \qquad | \qquad (s_n)$$
defined in
$$(3') \qquad | \qquad (3) \qquad (3)$$
Prop. 3. (Transitive property). If:
$$\sigma = \lambda, \ \sigma' = \lambda' \text{ and } \lambda = \lambda' \qquad | \qquad \lambda = \sigma, \ \lambda' = \sigma' \text{ and } \sigma = \sigma'$$
also
$$\sigma = \sigma' \qquad | \qquad \lambda = \lambda'$$
Proof. Being, by assumption
$$s_1 = s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \text{ es}_{\lambda} \qquad | \qquad s_n = s_{\sigma'} \sum_{n=0}^{\infty} a_n \text{ for } \sum_{n=0}^{\infty} a_n \in S_{\sigma} \text{ it results}$$

$$S_{\sigma} = S_{\sigma'} \text{ and } \sigma \sum_{n=0}^{\infty} a_n = \sigma' \sum_{n=0}^{\infty} a_n \text{ for } \sum_{n=0}^{\infty} a_n \in S_{\sigma} \qquad | \qquad s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n = \lambda' \lim s_n \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n \in s_{\lambda'} \text{ for } \{s_n\} \in s_{\lambda'} \text{ and } \lambda \lim s_n \in s_{\lambda'} \text{ for } \{s_n\} \in s_{\lambda'} \text{ for$$

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 $\sigma \subset \sigma' \ (or \ \sigma \subset \sigma' \ but \ \sigma \neq \sigma')$ $\lambda \subset \lambda' \ (or \ \lambda \subset \lambda' \ but \ \lambda \neq \lambda')$

 $\lambda = \sigma$ and $\lambda' = \sigma'$

 $\sigma = \lambda$ and $\sigma' = \lambda'$,

we have

Proof. Being

 $s_{\lambda} \subset s_{\lambda'} \ (or \ s_{\lambda} \subset s_{\lambda'} \ \text{but} \ s_{\lambda} \neq s_{\lambda'}), \qquad | \qquad S_{\sigma} \subset S_{\sigma'} \ (or \ S_{\sigma} \subset S_{\sigma'} \ \text{but} \ S_{\sigma} \neq S_{\sigma'}),$

it results

 $S_{\sigma} \subset S_{\sigma'}$ (or $S_{\sigma} \subset S_{\sigma'}$ but $S_{\sigma} \neq S_{\sigma'}$) $S_{\lambda} \subset S_{\lambda'}$ (or $S_{\lambda} \subset S_{\lambda'}$ but $S_{\lambda} \neq S_{\lambda'}$)

after

(1) and (2) (1') and (2')

Formals laws

Additive law

a) If $\{s_n\} \in \mathbf{s_\lambda} \text{ and } \{t_n\} \in \mathbf{s_\lambda},$ $\sum_{n=0}^{\infty} a_n \in \mathbf{S}_{\sigma} \text{ and } \sum_{n=0}^{\infty} b_n \in \mathbf{S}_{\sigma}$

also

 $\{s_n + t_n\} \in S_{\lambda}$ $\sum_{n=0}^{\infty} (a_n + h_n) \in S_{\sigma}$

and we have

 $\lambda \lim (s_n + t_n) = \lambda \lim s_n + \lambda \lim t_n$ $\sigma \sum_{n=0}^{\infty} (a_n + b_n) = \sigma \sum_{n=0}^{\infty} a_n + \sigma \sum_{n=0}^{\infty} b_n$

Homogeneity

b) If | b') If

 $\{s_n\} \in \mathbf{S}_{\lambda},$ $\sum_{n=0}^{\infty} a_n \in \mathbf{S}_{\sigma},$

also

 $\{c \ s_n\} \ \epsilon \ s_{\lambda} \qquad \sum_{n=0}^{\infty} c \ a_n \ \epsilon \ s_{\sigma}$

and it is

 $\lim (c s_n) = c. \lambda \lim s_n$ $\sigma \sum_{n=0}^{\infty} (c a_n) = c. \sigma \sum_{n=0}^{\infty} a_n$

whichever the constant c may be.

Initial term

c) If $\{s_n\} \in \mathbf{s}_{\lambda} \text{ and } \{s_{n+1}\} \in \mathbf{s}_{\lambda}, \qquad \sum_{n=0}^{\infty} a_n \in \mathbf{S}_{\sigma} \text{ and } \sum_{n=1}^{\infty} a_n \in \mathbf{S}_{\sigma},$

being

 $s_0 = o, a_0 = o$

we have

$$\lambda \lim s_n = \lambda \lim s_{n+1}$$

$$\sigma \sum_{n=0}^{\infty} a_n = \sigma \sum_{n=1}^{\infty} a_n$$

Product law

d) If

$$\{s_n\} \in s_{\lambda}, \{t_n\} \in s_{\lambda} \text{ and } \{z_n\} \in s_{\lambda}$$

 $\{s_n\} \in \mathbf{S}_{\lambda}, \{t_n\} \in \mathbf{S}_{\lambda} \text{ and } \{z_n\} \in \mathbf{S}_{\lambda}$ $\sum_{n=0}^{\infty} a_n \in \mathbf{S}_{\sigma}, \sum_{n=0}^{\infty} b_n \in \mathbf{S}_{\sigma} \text{ and } \sum_{n=0}^{\infty} c_n \in \mathbf{S}_{\sigma}$

being

$$z_n = s_n t_n \text{ for } n = 0, 1, 2, \dots$$

 $z_n = s_n t_n \text{ for } n = 0, 1, 2, \ldots,$ $c_n = a_n * b_n = \sum_{\nu=0}^n a_{\nu} b_{n-\nu} \text{ for } n = 0, 1, 2, \ldots,$

we have

$$\lambda \lim z_n = \lambda \lim s_n$$
. $\lambda \lim t_n$

$$\sigma \sum_{n=0}^{\infty} c_n = \sigma \sum_{n=0}^{\infty} a_n. \ \sigma \sum_{n=0}^{\infty} b_n$$

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma}, \sum_{n=0}^{\infty} b_n \in S_{\sigma} \text{ and } \sum_{n=0}^{\infty} c'_n \in S_{\sigma},$$

$$c'_0 = o$$
 and $c'_{n+1} = a_n * b_n$ for $n = 0, 1, 2, ...,$

$$\sigma \sum_{n=0}^{\infty} c'_n = \sigma \sum_{n=0}^{\infty} a_n. \ \sigma \sum_{n=0}^{\infty} b_n$$

Abel law

e₁) If

$$\{s_n\} \in S_{\lambda},$$

 $\sum_{n=0}^{\infty} a_n \in S_{\sigma},$ (4')(4)

also

$$\{n, x^n\} \in S_\lambda$$
 (

$$\sum_{n=0}^{\infty} a_n x^n \in \mathbf{S}_{\mathbf{d}} \tag{5'}$$

$$\left\{\sum_{\nu=0}^{n} s_{\nu} x^{\nu}\right\} \epsilon s_{\lambda}$$

$$\sum_{n=0}^{\infty} a_n x^n \in S_{\sigma}$$

$$\sum_{n=0}^{\infty} (1 * a_n) x^n \in S_{\sigma}$$
(6')

and we have

$$\lambda \lim (s_n x^n) = 0$$
 $\sigma \sum_{n=0}^{\infty} a_n x^n = (x-1) \cdot \sigma \sum_{n=0}^{\infty} (1 * a_n) x^n$ (7')

for any $x \in [0, 1]$

e₂) If

$$\{s_n\} \in S_{\lambda}$$

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma}$$

and $x \in]0$, 1[, we have

$$\lim_{x\to 1-} \left[(1-x) \cdot \lambda \lim_{\nu=0} \sum_{n=0}^{n} s_{\nu} x^{\nu} \right] = \lambda \lim_{n\to 1-} s_{n} (8)$$

$$\lim_{x\to 1-} \sigma \sum_{n=0}^{\infty} a_{n} x^{n} = \sigma \sum_{n=0}^{\infty} a_{n}$$

f) The limit

(8')

$$\sigma \sum_{n=0}^{\infty} a_n$$

is a linear continuous functional in the space

with a topology.

Remark. - When c'is verified, d' and d', are evidently equivalent.

Through the laws

a and b, the law c

a' and b', the law c'

can be generalized for any initial term as follows:

 c_1) If

$$\{s_n\} \in s_{\lambda} \text{ and } \{s_{n+1}\} \in s_{\lambda},$$

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma} \text{ and } \sum_{n=1}^{\infty} a_n \in S_{\sigma},$$

we have

$$\lambda \lim s_n = \lambda \lim s_{n+1}$$

$$\sum_{n=0}^{\infty} a_n = a_0 + \sum_{n=1}^{\infty} a_n$$

whichever the initial term may be.

Proof. Writing

$$s'_0 = 0$$
 and $s'_n = s_n$ for $n = 1, 2, ...$ $a'_0 = 0$ and $a'_n = a_n$ for $n = 1, 2, ...$ $s''_0 = s_0$ and $s''_n = 0$ for $n = 1, 2, ...$ $a''_0 = a_0$ and $a''_n = 0$ for $n = 1, 2, ...$

we have evidently

$$s_n = s'_n + s''_n \text{ for } n = 0, 1, 2, ...;$$
 $a_n = a'_n + a''_n \text{ for } n = 0, 1, 2, ...$

$$a''_0 = a_0$$
 and $a''_n = 0$ for $n = 1, 2, ...$

but us

$$\{s_n''\}\in S_1, \subset S_2, \lambda \lim s''_n = \lim s''_n = 0$$

$$a_n = a'_n + a''_n$$
 for $n = 0, 1, 2, ...$

 $\{s_n''\} \in s_{\lambda_0} \subset s_{\lambda}, \ \lambda \lim s''_n = \lim_{n \to \infty} s''_n = 0$ $\sum_{n=0}^{\infty} a_n'' \in S_{\sigma_0} \subset S_{\sigma}, \quad \lambda \sum_{n=0}^{\infty} s_n'' = \sum_{n=0}^{\infty} s_n'' = a_0$

it follows from

a and b,

a' and b',

that also

$$\{s'_n\} \in S_{\lambda};$$

$$\sum_{n=0}^{\infty} a'_n \in S_{\sigma};$$

wherefrom, being nul its initial term, it follows according to

c:
$$\lambda \lim s_n = \lambda \lim s'_n + \lambda \lim s''_n = \lambda \lim s_{n+1} +$$

$$c': \sigma \sum_{n=0}^{\infty} a_n = \sigma \sum_{n=0}^{\infty} a'_n + \sigma \sum_{n=0}^{\infty} a''_n = \sigma \sum_{n=1}^{\infty} a'_n +$$

$$+ \lambda_0 \lim s''_n = \lambda \lim s_{n+1}$$

$$+ \sigma_0 \sum_{n=0}^{\infty} a''_n = \sigma \sum_{n=1}^{\infty} a_n + a_0$$

From the laws

$$e_1$$
 and e_2 , e'_1 and e'_2 ,

can be deduced the following one, which may replace both in the applications.

e) If
$$\begin{cases} s_n \} \in s_{\lambda} \\ \end{cases} \qquad \qquad \begin{cases} e' \text{) If} \\ \sum_{n=0}^{\infty} a_n \in S_{\sigma} \end{cases}$$

it is also

$$\left\{\sum_{\nu=0}^{n} s_{\nu} x^{\nu}\right\} \epsilon s_{\lambda} \qquad \sum_{n=0}^{\infty} a_{n} x^{n} \epsilon S_{\sigma}$$

for every $x \in]0, 1[$, and we have

$$\lim_{x\to 1^-} \left[(1-x) \cdot \lambda \lim \sum_{\nu=0}^n s_{\nu} x^{\nu} \right] = \lambda \lim s_n \qquad \lim_{x\to 1^-} \sigma \sum_{n=0}^{\infty} a_n x^n = \sigma \sum_{n=0}^{\infty} a_n$$

Proof. The assumption is the same as in

$$e_1$$
 and e_2 | e'_1 and e'_2

and the conclusions are

(6) of
$$e_1$$
 and (8) of e_2 | (6') of e'_1 and (8') of e'_2

Whichever may be the formal laws taken in Def. 1 of the algorithms, the laws a, b, c, e_1, e_2 and f | a', b', c', e'_1, e'_2 and f'

are surely transformed in the

$$a', b', c', e'_1, e'_2$$
 and f'
 a, b, c, e_1, e_2 and f

in the equivalence (Def. 5) of

with another

as expressed in the following proposition

Prop. 5. If

a convergence algorithm
$$\lambda$$
 | an addition algorithm σ

verifies the formal laws a, b, c, e_1 , e_2 and fa', b', c', ε'_1 , ε'_2 and f'with a certain topology T, T_{σ} every addition algorithm o convergence algorithm h equivalent (Def. 5) to the σ , $\lambda = \sigma$, λ , $\sigma = \lambda$, verifies the homolous laws a', b', c', e'_1 , e'_2 and f' a, b, c, e_1, e_2 and fwith the topology T_{λ} T_{σ} , obtained by taking as neighborhood of a series $\sum_{n=0}^{\infty} a_n^{(0)} \in S_{\sigma}$ in S_{σ} , $\{s_n^{(0)}\}\ \epsilon\ s_{\lambda}\ in\ s_{\lambda}$ sequence the set of the series $\sum_{n=0}^{\infty} a_n \in S_{\sigma}$ sequences $\{s_n\} \in S_{\lambda}$ whose series $\sum_{n=0}^{\infty} a_n$ sequences $\{s_n\}$ given by (3) belongs to the neighborhood in after the topology T_{λ} with which is verified ť, of the sequence series obtained by applying

(3')

(3)

for the sequence

$$a_n = a_n^{(0)} \qquad \qquad s_n = s_n^{(0)}$$

for n = 0, 1, 2, ...

Proof. Being

$$s_0 \subset s_\lambda$$
 and $s_0 \neq s_\lambda$, $s_0 \subset s_0$ and $s_0 \neq s_0$,

it results from

$$(1): S_0 \subset S_\sigma \text{ and } S_0 \neq S_\sigma \qquad \qquad | \qquad \qquad (1'): s_\sigma \subset s_\lambda \text{ and } s_0 \neq s_\lambda$$

Furthermore, if

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma}, \qquad \{s_n\} \in S_{\lambda}$$

i. e., after

(1),
$$\{s_n\} \in s_\lambda$$

$$(1'), \sum_{n=0}^{\infty} a_n \in S_{\sigma}$$

we deduce from

being

an extension of

$$\sigma \sum_{n=0}^{\infty} a_n = \lambda \lim s_n = \lambda_0 \lim s_n = \lim_{n \to \infty} s_n = \begin{cases} s_n = 0 \\ s_n = 0 \end{cases}$$

$$\lambda \lim s_n = \sigma \sum_{n=0}^{\infty} a_n = \sigma_0 \sum_{n=0}^{\infty} a_n = \sum_{n=0}^{\infty} a_n = s_n = s$$

that is

is an extensión of

$$\sigma_0$$

Now, we are going to proof the formal laws

a') Let be
$$\sum_{n=0}^{\infty} a_n \in S_{\sigma} \text{ and } \sum_{n=0}^{\infty} b_n \in S_{\sigma},$$
 and $\{s_n\} \in S_{\lambda} \text{ and } \{t_n\} \in S_{\lambda},$

i. e., after

(1),
$$\{s_n\}$$
 ϵ s_{λ} and $\{t_n\}$ ϵ s_{λ} (1'), $\sum_{n=0}^{\infty} a_n \epsilon S_{\sigma}$ and $\sum_{n=0}^{\infty} b_n \epsilon S_{\sigma}$

being

$$s_n = \sum_{\nu=0}^n a_{\nu}$$
 and $t_n = \sum_{\nu=0}^n b_{\nu}$ for $n = 0, 1, 2, ...$ $a_0 = s_0, \ a_n = \Delta s_{n-1}$ and $b_0 = t_0, \ b_n = \Delta t_{n-1}$ for $n = 1, 2, ...$

Since

$$\sum_{v=0}^{n} (a_v + b_v) = s_n + t_n \text{ for } n = 0, 1, 2, \dots$$

we have, after

(2) and a: $\sigma \sum_{n=0}^{\infty} (a_n + b_n) = \lambda \lim (s_n + t_n) = \lambda \lim s_n + \lim (s_n + t_n) = \sigma \sum_{n=0}^{\infty} (a_n + b_n) = \sigma \sum_{n=0}^{\infty} a_n + \sigma \sum_{n=0}^{\infty} b_n + \lambda \lim t_n = \sigma \sum_{n=0}^{\infty} a_n + \sigma \sum_{n=0}^{\infty} b_n = \lambda \lim s_n + \lambda \lim t_n$

which is the

 $\begin{array}{c|c} a') & & & \\ b') \cdot \text{If} & & & \\ \sum_{n=0}^{\infty} a_n \in S_{\sigma}, & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & & \\ & &$

i. e., after

(1),
$$\{s_n\}$$
 ϵ s_{λ} $(1')$, $\sum_{n=0}^{\infty} a_n \epsilon S_{\sigma}$

also, after

b,
$$b\{c \mid s_n\} \in S_\lambda$$
 b', $\sum_{n=0}^{\infty} c \mid a_n \in S_\sigma$

whichever the constant c may be, that is, after

(1),
$$\sum_{n=0}^{\infty} c \ a_n \in S_{\sigma}$$
 (1'), $\{c \ s_n\} \in S_{\lambda}$

and besides, we have, after

b:
$$\sigma \sum_{n=0}^{\infty} (c \ a_n) = \lambda \lim_{n=0}^{\infty} (c \ s_n) = c. \ \lambda \lim_{n=0}^{\infty} s_n = c. \ \sigma \sum_{n=0}^{\infty} a_n = c. \ \sigma \sum_{n=0}^{\infty} a_n = c. \ \lambda \lim_{n=0}^{\infty} s_n = c. \ \lambda \lim_{n=0}$$

which is just

c') Let be

be c) Let be
$$\sum_{n=0}^{\infty} a_n \in S_{\sigma} \text{ and } \sum_{n=1}^{\infty} a_n \in S_{\sigma}$$
 (9)
$$\{s_n\} \in S_{\lambda} \text{ and } \{s_{n+1}\} \in S_{\lambda}$$

being

$$a_0 =$$

By according with

(3):
$$s_n = \sum_{\nu=0}^n a_{\nu}$$
 for $n = 0, 1, ...$ (3'): $a_0 = o, a_n = \Delta s_{n-1}$ for $n = 1, 2, ...$ $s'_n = \sum_{\nu=1}^{n+1} a_{\nu}$ for $n = 0, 1, ...$ $a'_0 = s_1, a_n = \Delta s_n$ for $n = 1, 2, ...$

(3'):
$$a_0 = o$$
, $a_n = \Delta s_{n-1}$ for $n = 1, 2, ...$

$$a'_0 = s_1, \ a_n = \Delta \ s_n \ \text{for} \ n = 1, 2, \dots$$

we have, after

(1) and (9),
$$\{s_n\} \in s_{\lambda}$$
 and $\{s'_n\} \in s_{\lambda}$

(1) and (9),
$$\{s_n\} \in s_{\lambda}$$
 and $\{s'_n\} \in s_{\lambda}$ (1') and (9'), $\sum_{n=0}^{\infty} a_n \in s_{\sigma}$ and $\sum_{n=0}^{\infty} a'_n \in s_{\sigma}$

But it is abviously

$$s'_n = s_{n+1}$$

$$a'_n = a_{n+1}$$

for $n = 0, 1, 2, \ldots$; whence it results by successive applying of:

(2), c and (2):

$$\sigma \sum_{n=0}^{\infty} a_n = \lambda \lim s_n = \lambda \lim s_{n+1} = \lambda \lim s'_n = \lambda \lim s_n = \sigma \sum_{n=0}^{\infty} a_n = \sigma \sum_{n=0}^{\infty} a_n = \sigma \sum_{n=0}^{\infty} a'_n = \sigma \sum_{n=$$

$$\lambda \lim s_n = \sigma$$

$$\sigma \sum_{n=0}^{\infty} a_n = \lambda \lim s_n = \lambda \lim s_{n+1} = \lambda \lim s'_n =$$

$$\lambda \lim s_n = \sigma \sum_{n=0} a_n = \sigma \sum_{n=1} a_n = \sigma \sum_{n=0} a'_n =$$

$$= \sigma \sum_{n=1}^{\infty} a_n$$

$$= \lambda \lim s_{n+1}$$

e'1) Let be

$$\{s_n\} \in S_{\lambda}$$

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma}$$

The conclusion

(6) of
$$e_1$$
,

is exactly the

after

(1) and (3)

From Abel's identity, or partial addition:

$$\sum_{\nu=0}^{n+1} a_{\nu} x^{\nu} = (1-x) \sum_{\nu=0}^{n} s_{\nu} x^{\nu} + s_{n+1} x^{n+1}, \qquad s_{n} x^{n} = \sum_{\nu=0}^{n} a_{\nu} x^{\nu} - (1-x) \sum_{\nu=0}^{n-1} (1*a_{\nu}) x^{\nu},$$

(10')

and through a new application of

$$(6)$$
, b , (5) , b , a and (1) ,

(6'), (1'), b, (5'), (1'), a and b,

it follows

$$\sum_{n=0}^{\infty} a_n \, x^n \in \mathbf{S}_{\sigma} \qquad \{s_n \, x^n\} \in \mathbf{S}_{\lambda}$$

which is

(5') of
$$e'_1$$
;

(5) of
$$e_1$$
:

and besides, by virtue of

(2) and (7) of
$$e_1$$
,

(2') and (7') of
$$e'_1$$

we have

$$\sigma \sum_{n=0}^{\infty} a_n \ x^n = (1-x) \ \lambda \lim_{\nu=0}^{n} s_{\nu} \ x^{\nu} + \lambda \lim_{\nu=0}^{n} (s_n x^n) = \lambda \lim_{\nu=0}^{n} \sum_{n=0}^{n} a_n \ x^{\nu} - (1-x) \cdot \lambda \lim_{\nu=0}^{n} \sum_{n=0}^{n} (1 * a_n) \ x^{\nu} = (1-x) \cdot \lambda \lim_{\nu=0}^{n} \sum_{n=0}^{\infty} (1 * a_n) \ x^{\nu} = (1-x) \cdot \sigma \sum_{n=0}^{\infty} s_n x^n = (1-x) \cdot \sigma \sum_{n=0}^{\infty} (1 * a_n) x^n = 0$$

$$(10)$$

which is just

(7') of
$$e'_1$$

(7) of
$$e_1$$
.

e'ı) From

(10), (8) of
$$e_2$$
 and (2)

(10), (8) of
$$e_2$$
 and (2), (10'), (2'), (8') of e'_2 and (2')

it results

$$\lim_{x \to 1^-} \int_{n=0}^{\infty} a_n x^n = \lim_{x \to 1^-} [(1-x) \cdot \lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{x \to 1^-} [(1-x) \cdot \lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 1^-} [(1-x) \cdot \lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [(1-x) \cdot \lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x^n] = \lim_{v \to 0} [\lambda \lim_{v \to 0} \sum_{n=0}^{n} s_n x$$

$$\lim_{x\to 1^{-}} [(1-x). \lambda \lim_{v\to 0} \sum_{v=0}^{n} s_{v} x^{v}] = \lim_{v\to 1^{-}} [(1-x).$$

$$\lambda \lim_{x\to 1^{-}} \sum_{v=0}^{n} (1*a_{v}) x^{v}] = \lim_{v\to 1^{-}} [\lambda \lim_{v\to 1^{-}} \sum_{v=0}^{n} a_{v} x^{v}] =$$

$$= \lim_{x \to 1^{-}} \sigma \sum_{n=0}^{\infty} a_n x^n = \sigma \sum_{n=0}^{\infty} a_n = \lambda \lim s_n$$

which is the conclusion

(8') of
$$e'_2$$

(8) of
$$e_2$$

f') After

f,

to every ε-neighborhood of the number

$$\sigma \sum_{n=0}^{\infty} a_n^{(0)} = \lambda \lim s_n^{(0)}$$

$$\lambda \lim s_n^{(0)} = \sigma \sum_{n=0}^{\infty} a_n^{(0)}$$

it corresponds a $\delta\text{-neighborhood}$ of

$$\{s_n^{(0)}\} \qquad \sum_{n=0}^{\infty} a_n^{(0)}$$

in the topology

$$T_{\lambda}$$

such that if

$$\{s_n\} \in \delta, \quad \lambda \lim s_n \in \varepsilon; \quad (11) \mid \sum_{n=0}^{\infty} a_n \in \delta, \quad \sigma \sum_{n=0}^{\infty} a_n \in \varepsilon; \quad (11')$$

and to this one, another δ' of

$$\sum_{n=0}^{\infty} a_n^{(0)}$$
 $\{s_n^{(0)}\}$

in the topology

$$T_{\lambda}$$

such that if

$$\sum_{n=0}^{\infty} a_n \in \delta', \qquad \{s_n\} \in \delta'$$

we have by definition

$$\{s_n\} \in \delta;$$

$$\sum_{n=0}^{\infty} a_n \in \delta;$$

therefore, after

we have

$$\sigma \sum_{n=0}^{\infty} a_n = \lambda \lim s_n \in \varepsilon$$

$$\lambda \lim s_n = \sigma \sum_{n=0}^{\infty} a_n \in \varepsilon$$

That is to say, to every ϵ -neighborhood of

$$\sigma \sum_{n=0}^{\infty} a_n^{(0)} \qquad \qquad \lambda \lim s_n^{(0)}$$

it corresponds another δ' of

$$\sum_{n=0}^{\infty} a_n^{(0)}$$
 $\{s_n^{(0)}\}$

in the topology

$$T_{\lambda}$$

such that if

$$\sum_{n=0}^{\infty} a_n \in \delta' \qquad (s_n) \in \delta'$$

we have

$$\sigma \sum_{n=0}^{\infty} a_n \in \varepsilon; \qquad \lambda \lim s_n \in s;$$

which means that the functional

$$\sigma \sum_{n=0}^{\infty} a_n \qquad \lambda \lim s_n$$

is continuous in

with the topology

$$T_{\sigma}$$

transformed from the topology

$$T_{\lambda}$$
 | T_{σ}

Remark. 1. In this Prop. 5 we have to do without the product laws, which are not deducible from an algorithm to the other one, since obviously

$$\sum_{\nu=0}^{n} a_{\nu}. \sum_{\nu=0}^{n} b_{\nu} \neq \sum_{\nu=0}^{n} (a_{\nu} * b_{\nu}),$$

and this is just the cause of troubles involved in the product law for summation algorithms, including those for convergent series.

2. In stead of these product laws and as auxiliaires for the converse deduction of Abel's laws

$$e$$
 or e_1

we have introduced the converse laws e_1 and e'_1 , whose last conclusions

are two particular cases of the said product laws, since we have

$$\lambda \lim (s_n x^n) = \lambda \lim s_n. \ \lambda \lim x^n = 0$$

$$= \lambda \lim s_n. \ \lambda_0 \lim x^n = 0$$

$$= \sigma \sum_{n=0}^{\infty} (1 * a_n) x^n = \sigma \sum_{n=0}^{\infty} (x^n * a_n x^n) = 0$$

$$= \sigma \sum_{n=0}^{\infty} x^n. \ \sigma \sum_{n=0}^{\infty} a_n x^n = \sigma_0 \sum_{n=0}^{\infty} x^n. \ \sigma \sum_{n=0}^{\infty} a_n x^n = 0$$

$$= \frac{1}{1-x}. \ \sigma \sum_{n=0}^{\infty} a_n x^n$$

2. Linear algorithms

It is convenient to fix some terminology through the following definitions:

Def. 1. We will call linear algorithms of

convergence,
$$\lambda$$
 addition, μ

those defined by means of a matrix, or double sequence

$$\{\lambda_{r,t}\}$$
 | $\{\mu_{r,t}\}$

with two indexes r = 0, 1, 2, ... and t = 0, 1, 2, ... as in

and linear algorithms of

convergence,
$$\lambda$$
 addition, μ

with continuous parameter, t, or more shortly continuous algorithms, those defined by means of a factors sequence of

$$\{\lambda_r(t)\}$$
 - convergence $\{\mu_r(t)\}$ - addition

with an r = 0, 1, 2, ... and variable t within a set E with the adherence point t_0 , as in [12, 193-197], where it was analyzed in detail the relation between them.

Permanent linear algorithms are evidently distributive algorithms (Def. 1, no. 1).

Def. 2. We call transforms of a

sequence
$$\{a_n\}$$
 (1) $\sup_{n=0}^{\infty} a_n$ (1')

by means of a matrix

$$\lambda_{t,n}$$
 $\{\mu_{t,n}\}$

or by means of a sequence of functions

$$\{\lambda_n(t)\}$$
 $\{\mu_n(t)\}$

to the sequence

$$\left\{ \sum_{n=0}^{\infty} \lambda_{t,n} \ \alpha_n \right\} \tag{2'}$$

or the function of t, defined by the series

$$\sum_{n=0}^{\infty} \lambda_n (t) \alpha_n \qquad (3) \qquad \sum_{n=0}^{\infty} \mu_n (t) \alpha_n \qquad (3')$$

respectively.

Prop. 1. If the

form a vectorial space, their transforms

do it also.

Proof. Because the transformation is linear.

Def. 3. We will say that a linear algorithm of

convergence,
$$\lambda$$
 | addition, μ

verifies the principle of identity, when the only

sequence
$$(a_n)$$
 ϵ s_{λ} series $\sum_{n=0}^{\infty} a_n \epsilon s_{\mu}$

having all its transforms null

$$(2) (2')$$

for t = 0, 1, 2, ... or the

for $t \in E$, is the identically zero one

$$a_n = 0$$
 for $n = 0, 1, 2, ...$

and when furthermore, in the second case, the set E is compact and the transforms

$$(2) (2'$$

of every

sequence
$$\{s_n\}$$
 ϵ s_{λ} series $\sum_{n=0}^{\infty} a_n \epsilon S_{\mu}$

are continuous functions in $E-t_0 = E \cap \int_0^1 t_0$.

Prop. 2. If a linear algorithm verifies the identity principles, the set

$$\mathbf{S}_{\mu}$$

and that of the transforms correspond one-to-one each other; the vectorial space of the transforms

2) or
$$(3)$$
 $(2')$ or $(3')$

is a vectorial subspace \mathbf{c}_1 or \mathbf{c}_1 , respectively for the space \mathbf{c} of the convergent sequences or for the space \mathbf{c} of the continuous functions within the compact set E; and we may define a metric in the set

$$\mathsf{S}_{\pmb{\mu}}$$

in such a way that the latter resulte isometric with the subspace $c_1 \subset c$ or with the $C_1 \subset C$, adopting as norm for every

sequence
$$\{s_n\}$$
 ϵ s_2 , $series \sum_{n=0}^{\infty} a_n \epsilon S_{\sigma}$

that of its transform in c_1 or c_1 .

Proof. The correspondence between.

$$\mathbf{s}_{\lambda}$$
 and \mathbf{c}_1 or \mathbf{c}_1 \mathbf{s}_{μ} and \mathbf{c}_1 or \mathbf{c}_1

is obviously univocal; and one-to-one becase of the identity principle.

We have $c_1 \subset c$, because there exists

$$\lambda \lim s_n = \lim_{t \to \infty} \sum_{n=0}^{\infty} \lambda_t, n \ s_n < \infty \qquad (5) \qquad \sigma \sum_{n=0}^{\infty} a_n = \lim_{t \to \infty} \sum_{n=0}^{\infty} \mu_t, n \ a_n < \infty \qquad (5')$$

and $C_1 \subset C$ because existing

$$\lambda \lim s_n = \lim_{t \to t_0} \sum_{n=0}^{\infty} \lambda_n(t) \, s_n < \infty, \ t \in E, (6)$$

$$\sigma \sum_{n=0}^{\infty} a_n = \lim_{t \to t_0} \sum_{n=0}^{\infty} \mu_n(t) \, a_n, \ t \in E \quad (6')$$

the transform is furthermore continuous in E- t_0 ; therefore this transform is a cotinous function in the whole set E, adopting this limit as value at t_0 .

Finally, the norm has identical property by virtue of the identity principle; and the spaces result obviously isometric, since they are vectorial spaces with equal norms at homologous points.

Def. 4. The metric defined in Prop. 2 and the topology determined in

will be called specific for the algorithm (*).

Prop. 3. If an algorithm verifies the identity principle (Def. 3) it satisfies the continuity law

$$f$$
 | f'

with its specific topology.

Proof. The limit

is, of course, a continuous functional in c_1 or c_1 with the metric of c or c, respectively, and, therefore, in

$$\mathsf{S}_{\mu}$$

with the specific metric of the algorithm, because then, both spaces are isometric; and therefore it shall be continuous with the specific topology for the algorithm (*).

The definitions of reversible and perfect convergence algorithms [1, 90] could be generalized for continuous algorithms. But the extension of fundamental theorems requires a detailed study, specially the [1, Lemma 4, 93], based on the [1, Theor. 10, 47] which we are trying to avoid here, through the algorithms with matrix subordinated to a continuous algorithm.

Def. 5. Being given a continuous linear algorithm of

convergence,
$$\lambda$$
 addition, μ

with the factors

$$\{\lambda_n(t)\},\qquad \qquad \{\mu_n(t)\},$$

we call subordinate algorithm through a sequence $\{r_t\}$ ϵ E with limit t_0 and it will be designed sometimes by

$$\lambda (r_t), \qquad \qquad | \qquad \qquad \mu (r_t),$$

to the algorithm of

with the matrix

$$\{\lambda_{\mathbf{n}}(r_t)\} \qquad \qquad \{\mu_{\mathbf{n}}(r_t)\}$$

for t = 0, 1, 2, ... and n = 0, 1, 2, ...

λ

^(*) Somewhere else we will analize the relation between this metric and the topology defined by Erdös and Piranian [6, 139-148] by means of the Toeplitz matrices in the set of bounded sequences.

be a continuous algorithm of

In order that there exist

$$\lambda \lim s_n,$$
 (7)
$$\mu \sum_{n=0}^{\infty} a_n,$$
 (7')

that is,

$$\{s_n\} \in S_{\lambda},$$

$$\sum_{n=0}^{\infty} a_n \in S_{\mu}$$

it is necessary and sufficient that there exist

$$\lambda(r_t) \lim s_n, \qquad (8) \qquad \mu(r_t) \sum_{n=0}^{\infty} a_n, \qquad (8')$$

that is,

$$\{s_n\} \in S_{\lambda(r_t)},$$

$$\sum_{n=0}^{\infty} a_n \in S_{\mu(r_0)}$$

for every sequence $\{r_t\}$ ϵ E with limit t_0 . Evidently:

result equal each other and

Using symbols

$$\mathbf{s}_{\lambda} = \bigcap \left(\mathbf{s}_{\lambda(r_t)}, \{ r_t \} \in E, \lim_{t \to \infty} r_t = t_0 \right) \quad (9)$$

$$\lambda \lim s_n = \lambda(r_t) \lim s_n \quad (10)$$

$$\mathbf{s}_{\mu} = \bigcap \left(\mathbf{s}_{\mu(r_t)}, \{ r_t \} \in E, \lim_{t \to \infty} r_t = t_0 \right) \quad (9')$$

$$\mu \sum_{n=0}^{\infty} a_n = \mu(r_t) \sum_{n=0}^{\infty} a_n \quad (10')$$

for every sequence $\{r_t\} \in E$.

In particular, if the compact set E, with the adherence point t_0 is the real positive semi-axis $[0, +\infty]$ and $t=+\infty$, the preceding conclusions remain valid, but limited to the increasing sequences $\{r_t\}$ with limit $+\infty$ and all their terms larger than any prefixed constant.

Proof. They are abvious consequences of the elemental properties of arithmetical and functional limits.

Cors. 1. Any continuous algorithm is weaker than any of oits subordinate ones, i. e.

$$\lambda \subset \lambda (r_t)$$
 $\mu \subset \mu (r_t)$

for any sequence $\{r_t\}$ ϵ E with limit t_0 .

Proof. An obvious consequence from

2. If a continuous algorithm is permanent, all its subordinates are permanent too.

Proof. Immediate consequence from permanence definition:

$$\lambda_0 \subset \lambda$$
, $\lambda_0 \lim s_n = \lambda \lim s_n$, $\mu_0 \subset \mu \ (\mu_0 = \sigma_0)$, $\mu_0 \sum_{n=0}^{\infty} a_n = \mu \sum_{n=0}^{\infty} a_n$

(Prop. 1, no. 1), Cor. 1 and

3. Moments method.

We are still giving immediate definitions and propositions, which will be used later.

summatory
$$\{v_n\}$$
 integral $\{\mu_n\}$

algorithm of moments

to a permanent method [13, 34] of summation, defining the sum of a series $\sum_{n=1}^{\infty} a_n$ as follows

$$\nu \sum_{n=0}^{\infty} a_n = \sum_{t=0}^{\infty} a_t A_t < \infty \qquad (1) \qquad \mu \sum_{n=0}^{\infty} a_n = \int_0^{\infty} a(t) A(t) dt < \infty \qquad (1')$$

where

 $\{a_t\}$ is the generatrix sequence with the $|a_t|$ is the continuous function in $[0, +\infty]$ moments

$$\nu_n = \sum_{t=0}^{\infty} a_t \ t^{(n)} < \infty$$

for $n = 0, 1, 2, \ldots$ and A_t the associated sequence formed by the values of the polvnomials

$$A_t = \sum_{n=0}^t a_n \frac{t^{(n)}}{n}$$
 for $t = 0, 1, 2, ...$ (2)

Prop. 1. Conversely, any

sequence $\{A_t\}$ is the associated

$$\mu_{n} = \int_{0}^{\infty} a(t) t^{n} dt < \infty$$

and A(t) the associated function obtained by analytical continuation along the positive real semi-axis of the function holomorph in 0 defined by the series

$$A(t) = \sum_{n=0}^{\infty} a_n \frac{t^n}{n!}$$
 (2')

holomorph function, A(t), in 0, analytical continuable along the whole positive real semiaxis is the associated

of a series $\sum_{n=0}^{\infty}$ with coefficients

$$a_0 = A_0 = \Delta^0 A_0$$
 and $a_n = \frac{\Delta^n A_0}{n!} \nu_n$ (3) $a_0 = A(0) = A(0) (0)$ and $a_n = \frac{A(n)(0)}{n!} \mu_n$ (3')

for $n=1, 2, \ldots$

Proof. Since we have evidently:

$$\sum_{n=0}^{t} a_n \frac{t^{(n)}}{v_n} = \sum_{n=0}^{t} \frac{t^{(n)}}{n!} \Delta^n A_0 = \sum_{n=0}^{t} \binom{t}{n} \Delta^n A_0 = \begin{bmatrix} \sum_{n=0}^{\infty} a_n \frac{t^n}{\mu_n} = \sum_{n=0}^{\infty} \frac{t^n}{n!} A^{(n)} & (0) = A(t) \\ = A_t \text{ for } t = 0, 1, 2, \dots \end{bmatrix}$$
 in a neighborhood of the origin and by

$$\sum_{n=0}^{\infty} a_n \frac{t^n}{\mu_n} = \sum_{n=0}^{\infty} \frac{t^n}{n!} A^{(n)} (0) = A(t)$$

in a neighborhood of the origin and by continuation along the whole positive real semiaxis, since, by assumption, A(t) is holomorph in 0 and continuable along this semi-

Prop. 2. In the set

 S_{μ}

of summable series with an algorithm of

summatory moments {v_n} integral moments {\mu_n}

it can be defined a metric by adopting as norm for every series

$$\sum_{n=0}^{\infty} a_n \in S_r \qquad \sum_{n=0}^{\infty} a_n \in S_{\mu}$$

that of the product of the associated by the generatrix in the vectorial subspace

$$\mathfrak{U}_0 \subset \mathfrak{U}$$
 (4')

formed by the said products, so that we have a space

 \mathbf{S}_{μ}

isometric with

 \mathbf{u}_0 \mathbf{U}_{0}

Proof. By virtue of formulae

(2) and (3), (2') and (3'),

there exists a one-to-one correspondence, without exception, between the set

of summable series

and the set

 \mathbf{U}_0

of the products

 $\{\alpha_t A_t\}$ α (t) A (t)

of the generatrix,

 $\{\alpha_t\}$ $\alpha(t)$

of the algorithm, by the associated in

 $\{A_t\}$, A(t), of every series

$$\sum_{n=0}^{\infty} a_n \in S_{\nu} \qquad \sum_{n=0}^{\infty} a_n \in S_{\mu};$$

and in this correspondence, the identically zero series:

$$\sum_{n=0}^{\infty} a_n \text{ with } a_n = 0 \text{ for } n = 0, 1, 2, \dots$$

and the identically zero associated

$$A_t=0$$
 for $t=0, 1, 2, \ldots$ | $A(t)=0$ for $0 \le t < +\infty$

become homologous.

This set

 u_0 U

is evidently a vectorial space, as well as the

 S_{ν}

But being convergent the

series (1), integrals (1')

the space

 u_0 | U_0

is contained in the

u | U

defined in

Then, we can adopt as norm for every series

S,

$$\sum_{n=0}^{\infty} a_n \in S_{\nu}$$

$$\sum_{n=0}^{\infty} a_n \in S_{\mu}$$

that of the homologous product

$$\{a_t \ A_t\} \ \epsilon \ u_0 \subset u$$
 $[a\ (t)\ A\ (t)] \ \epsilon \ U_0 \subset U$

in the space

This norm verifies the identical property, since, as we have said, the identically zero series and the identically zero product are homologous and so, both spaces

$$S_{\nu}$$
 and U_0 | S_{μ} and U_0

become isometrical, because they are vectorial spaces with equal norms at homologous points.

Def. 2. The metric defined in the Prop. 2 and the topology determined by this metric in

$$|$$
 S_{μ}

will be called specific for the algorithm

$$\mu$$

Prop. 3. The continuity law f' become true for the moments algorithm with the specific topology of the latter.

Proof. The value of the

being a continuous functional in the vectorial subspace

with the metric of the space

shall be also, according to Prop. 2, a continuous functional in

$$S_{\nu}$$
 |

with the metric and, therefore, with the specific topology of

Def. 3. An algorithm μ of integral moments being given (Def. 1) we will call subordinate algorithm to the μ through a sequence $\{r_t\}$ of positive numbers with $+\infty$ limit, and we will design it sometimes by μ (r_t) , to that one defining the sum of a series

$$\sum_{n=0}^{\infty} a_n$$
, as follows:

$$\mu (r_t) \sum_{n=0}^{\infty} a_n = \lim_{t \to \infty} \int_0^{r_t} a(v) A(v) dv$$

Prop. 4. In order to exist the sum:

$$\mu \sum_{n=0}^{\infty} a_n, \text{ that is, } \sum_{n=0}^{\infty} a_n \in S_{\mu},$$

it is necessary and sufficient that there exist the sums

$$\mu(r_t) \sum_{n=0}^{\infty} a_n$$
, that is, $\sum_{n=0}^{\infty} a_n \in S_{\mu(r_t)}$

for any sequence $\{r_i\}$ of positive numbers with $+\infty$ limit. Evidently all these sums are equal each other and to the sum μ .

In symbols:

$$\mathbf{S}\mu = \mathbf{\cap} (\mathbf{S}_{\mu(r_t)} \{r_t\} \in [0, +\infty], \lim_{t \to \infty} r_t = +\infty)$$

$$\mu \sum_{n=0}^{\infty} a_n = \mu(r_t) \sum_{n=0}^{\infty} a_n \text{ for any } \{r_t\} \in [0, +\infty] \text{ with } \lim_{t\to\infty} r_t = +\infty$$

Proof. Immediate consequence from elemental properties of limits.

Remark. The algorithm of

summatory integral

moments is linear with summation factors

$$\mu_{t,n} = \frac{1}{\nu_{n}} \sum_{\nu=0}^{t} \alpha_{\nu} \nu^{(n)} \text{ for } t = 0, 1, 2, \dots; \qquad \mu_{n}(t) = \frac{t}{\mu_{n}} \int_{0}^{t} x(\nu) \nu^{n} d\nu \text{ for } 0 = t \leq +\infty;$$

$$n = 0, 1, 2, \dots$$

$$n = 0, 1, 2, \dots$$

$$\mu_{n}(t) = \frac{t}{\mu_{n}} \int_{0}^{t} x(v) v^{n} dv \text{ for } 0 = t \leq +\infty;$$

$$n = 0, 1, 2, \dots$$

when the associated is an integral function; since then the series ((5'), no. 2) can be added together under the integral sign, because of the uniform convergence in $0 \le r \le t$

of
$$a(v) \sum_{n=0}^{\infty} a_n \frac{v^n}{\mu^n}$$

As, after

$$((3); no. 2),$$
 $((3'), no. 2),$

the algorithm satisfies evidently the edentity principle (Def. 3, no. 2), the foregoing Props. 2, 3 and 4, are included in the (Props. 2, 3 and 4, no. 2) respectively.

In particular, the metric or specific topology for the moment algorithm (Def. 2) coincides, then, in this case, with that one defined (Def. 4, 3, 2) for the linear algorithm, and the same happens with the definition of subordinate algorithms (Def. 3) which coincides with that other given by the linear ones (Def. 5, no. 2).

But, in the continuous case, the moments algorithm us more general because it does not requiere convergence of the associated series in the whole plan, but only the analytical continuation of the associated function along the positive real semiaxis. For this reason we have thought necessary to establish both directly for the moment algorithms.

4. The Euler's and Borel's associateds

Def. 1. As it is known the summation algorithm of

which will be for short designed by

$$E$$
 B

is a moments algorithm with the generatrix

$$a_t = 2^{-t-1}$$
 for $t = 0, 1, 2, ...$ $a(t) = e^{-t}$ for $0 \le t < +\infty$

of moments

$$n! = \sum_{t=n}^{\infty} 2^{-t-1} t^{(n)}$$
 $n! = \int_{0}^{\infty} e^{-t} t^{n} dt$

for n = 0, 1, 2, ...

The associated is, thus, the

sequence $\{A_t\}$ formed by the values of the polynomials

function A (t) obtained by analytical continuation along the real positive semi-axis of the holomorph function in 0, defined by the series

$$A_t = \sum_{n=0}^t a_n \frac{t^{(n)}}{n!}$$
 for $t = 0, 1, 2, ...$

$$A(t) = \sum_{t=0}^{\infty} a_n \frac{t^n}{n!}$$

Tais associated has, among other important properties, the following, which we are going to use:

Prop. 1. The asociated

Proof. Since

$$A_{t} = \begin{bmatrix} a_{0} + a_{1} \begin{pmatrix} t+1 \\ 1 \end{pmatrix} + \dots + a_{i} \begin{pmatrix} t+1 \\ i \end{pmatrix} + A'(t) = a_{1} + a_{2} \frac{t}{1!} + \dots + a_{n} \frac{t^{n-1}}{(n-1)!} + \dots \\
+ \dots + a_{t} \begin{pmatrix} t+1 \\ t \end{pmatrix} + a_{t+1} \begin{pmatrix} t+1 \\ t+1 \end{pmatrix} \end{bmatrix} -$$

$$- \begin{bmatrix} a_{0} + a_{1} \begin{pmatrix} t \\ 1 \end{pmatrix} + \dots + a_{t} \begin{pmatrix} t \\ i \end{pmatrix} + \dots + a_{t} \begin{pmatrix} t \\ t-1 \end{pmatrix} + \dots + a_{t} \begin{pmatrix} t \\ t \end{pmatrix} \end{bmatrix} =$$

$$= a_{1} + a_{2} \begin{pmatrix} t \\ 1 \end{pmatrix} + \dots + a_{t} \begin{pmatrix} t \\ i-1 \end{pmatrix} + \dots + a_{t+1} \begin{pmatrix} t \\ t \end{pmatrix}$$

$$+ a_{t} \begin{pmatrix} t \\ t-1 \end{pmatrix} + a_{t+1} \begin{pmatrix} t \\ t \end{pmatrix}$$

for
$$t = 1, 2, ...$$
 and $\Delta A_0 = a_1$

which is the associated of $\sum_{n=0}^{\infty} a_n$, q. e. d.

Prop. 2. If
$$\sum_{n=0}^{\infty} a'_n$$
, with $a'_0 = 0$ and $a'_{n+1} = a_n$ for $n = 0, 1, 2, ...$ (1) has the Euler associated A'_t with $A'_0 = 0$ and $A'_{t+1} = A_t$ for $t = 0, 1, 2, ...$ (2) the $\sum_{n=0}^{\infty} a''_n$ with

Prop. 2'. The series
$$\sum_{n=0}^{\infty} a'_n$$
 with $a'_0 = 0$ and $a'_{n+1} = a_n$ for $n = 0, 1, 2, ...$ (1') has the Borel associated
$$F(t) = \int_{-\infty}^{t} A(v) dv$$

$$a^{\prime\prime}_{n}=a^{\prime}_{n}+a^{\prime}_{n+1}$$
 for $n=0,1,2,...$ that is $a^{\prime\prime}_{0}=a_{0}$ and $a^{\prime\prime}_{n}=a_{n-1}+a_{n}$ for $n=1,2,...$ has the Euler associated $\{A_{t}\}$

Proof. The associated of $\sum_{n=1}^{\infty} a^{n}$ is $A''_{t} = a_{0} + (a_{0} + a_{1}) \binom{t}{1} + \dots + (a_{i-1} + a_{i}) \binom{t}{i} + \qquad F(t) = a_{0} t + a_{1} \frac{tr}{2!} + \dots + a_{n-1} \frac{t^{n}}{n!} + \dots = a_{n-1} \frac{$ $+ \ldots + (a_{t-1} + a_t) \binom{t}{t} = a_0 \left[1 + \binom{t}{1} \right] + \qquad = a'_0 + a'_1 \frac{t}{1!} + \ldots + a'_n \frac{t^n}{n!} + \ldots$ $+a_1\left[\binom{t}{1}+\binom{t}{2}\right]+\ldots+a_{i-1}\left[\binom{t}{i-1}+\binom{t}{i}\right]+$ for $0 \le t < +\infty$, which is the associated $+ \dots + a_{t-1} \left[{t \choose t-1} + {t \choose t} \right] + a_t {t \choose t} = a_1 {t+1 \choose 1} +$ of $\sum_{n=0}^{\infty} a_n \neq 0$ e. d. $+ a'_{2}\binom{t+1}{2} + \dots + a'_{i}\binom{t+1}{i} + \dots +$ $+ a't \binom{t+1}{1} + a't+1 \binom{t+1}{t+1} = A't+1 = At$

Prop. 3. If

$$\{A_t\}$$
 and $\{B_t\}$

for t = 0, 1, 2, ... q. e. d.

are the associated

of $\sum_{n=0}^{\infty} a_n$ and $\sum_{n=0}^{\infty} b_n$, respectively, and we set

$$c_n = a_n * b_n = \sum_{\nu=0}^n a_{\nu} b_{n-\nu} \text{ for } n = 0, 1, 2, \dots$$

the associated of $\sum_{n=0}^{\infty} c'_n$ with

$$c'_0 = 0$$
 and $c'_{n+1} = c_n$ for $n = 0, 1, 2, ...$ (3)

is the

sequence
$$\{C'_t\}$$

function C(t)

Proof. Since

$$F(t) = a_0 t + a_1 \frac{t^n}{2!} + \ldots + a_{n-1} \frac{t^n}{n!} + \ldots =$$

$$= a'_0 + a'_1 \frac{t}{1!} + \ldots + a'_n \frac{t^n}{n!} + \ldots$$

defined by

$$C'_0 = 0$$
 and $C'_{t+1} - C_t = A_t * B_t =$

$$= \sum_{v=0}^{t} A_v B_{t-v} \text{ for } t = 0, 1, 2, \dots (4)$$

Proof. By forming the table of differences, we obtain

$$A^{n} C'_{0} = A_{0} A^{n-1} B_{0} + A_{0} A_{0} A^{n-2} B_{0} + A_{0} A^{n-3} B_{0} + \dots + A^{n-3} A_{0} A^{n-2} B_{0} + A^{n-2} A_{0} A_{0} B_{0} + A^{n-1} A_{0} B_{0}$$
(5) for $n = 1, 2, \dots$

waich is easily proved by induction.

In fact, it is evidently

$$A C'_0 = A_0 B_0$$

and being, after (4)

$$C'_t = \sum_{n=0}^{t-1} A_n B_{t-1-n}$$
 for $t = 1, 2, ...$

we have successively

$$A C'_{t} = \sum_{v=0}^{t-1} A_{v}...1 B_{t-1-v} + A_{t} B_{0}$$

$$12 C'_{t} = \sum_{v=0}^{t-1} A_{v} ... 12 B_{t-1-v} \oplus A_{t} ... 1B_{0} \oplus ... 1A_{t} ... B_{0}$$

$$A^{3} C'_{t} = \sum_{r=0}^{t-1} A_{r}, A^{3} B_{t-1-r} + A_{t}, A^{2} B_{0} + A_{t}, A B_{0} + A^{2} A_{t}, B_{0}$$

 $A^n C_t' := \sum_{\nu=0}^{t-1} A_{\nu} A_{\nu} A^n B_{t-1-\nu} + A_t A_t A^{n-1} B_0 = 0$

 $+ .1 A_t. A^{n-2} B_0 + ...$

$$\begin{split} A^n \, C'_1 &= A_0, \, A^n \, B_0 \, + \, A_1, \, \Delta^{n-1} \, B_0 \, + \\ &+ \, \Delta \, A_1, \, A^{n-2} \, B_0 \, \oplus \, \Delta^2 \, A_1, \, \Delta^{n-3} \, B_0 \, \oplus \ldots \, + \\ &+ \, A^{n-3} A_1, \, \Delta^2 B_0 + A^{n-2} A_1, \, \Delta B_0 + A^{n-1} A_1, B_0 \end{split}$$

$$C(t) = A * B = \int_0^t A(v) B(t-v) dv$$

$$for 0 < t < +\infty$$
 (4')

Proof. It is deduced from [5, 11-12] by virtue of Prop. 2'; since if w(t) is the associated of $\sum_{n=0}^{\infty} c_n$, that of $\sum_{n=0}^{\infty} c'_0$ is $\int_0^t W(v) dv$ according to Prop. 2'; and by virtue of [5, 11-12] we have

$$\int_{0}^{t} W(v) dv = C(t)$$

whence, after (5):

$$\begin{split} \varDelta^{n+1} \ C'_0 &= \varDelta^n \ C'_1 - \varDelta^n \ C_0' = A_0. \ \varDelta^n \ B_0 \ + \\ &+ \varDelta \ A_0 \ \varDelta^{n-1} \ B_0 \ + \ \varDelta^2 \ A_0 \ \varDelta^{n-2} \ B_0 \ + \\ &+ \varDelta^3 \ A_0. \ \varDelta^{n-3} \ B_0 \ + \ \dots \ + \varDelta^{n-2} \ A_0. \ \varDelta^2 \ B_0 \ + \\ &+ \varDelta^{n-1} \ A_0. \ \varDelta \ B_0 \ + \ \varDelta^n \ A_0. \ B_0 \end{split}$$

which is just the (5) for n+1.

But being $\{A_t\}$ and $\{B_t\}$ the associateds of $\sum_{n=0}^{\infty} a_n$ and $\sum_{n=0}^{\infty} b_n$, respectively, we have ((3), no. 3):

 $a_n = \Delta^n A_0$ and $b_n = \Delta^n B_0$ for n = 0, 1, 2, ...; then, after (5), it is

$$A^{n} C'_{0} = a_{0} b_{n-1} + a_{1} b_{n-2} + \dots + a_{n-1} b_{0} =$$

$$= c_{n-1} = c'_{n}$$

for n = 1, 2, ...; and as furthermore

$$C'_0 = 0$$
 and $c'_0 = 0$

it results actually, after ((3), no. 3) that $\{C'_0\}$ is the associated of $\sum_{n=0}^{\infty} c'_n$, q. e. d.

Re.nark. For the series $\sum_{n=0}^{\infty} a_n$ with

$$l = \lim_{n \to \infty} |a_n|^1 I^n < \infty,$$

the conclusion of this proposition becomes easily from the identity:

$$\sum_{n=0}^{\infty} a_n z^{n+1} = \sum_{t=0}^{\infty} A_t \left(\frac{z}{z+1} \right)^{t+1}$$

which remains valid [7, 7 and 179] for |z| < 1/l. But although these are the only summable series E, the restriction is artificial and needless for the validity of the result.

Prop. 4. If $\{A_t\}$ and $\{B_t\}$ are the associated of $\sum_{n=0}^{\infty} a_n$ and $\sum_{n=0}^{\infty} b_n$, respectively, and

we set $c_n = a_n * b_n$ for n = 0, 1, 2, ... the associated of $\sum_{n=0}^{\infty} c''_n$ with

 $c''_0 = c_0$ and $c''_n = c_{n-1} + c_n$ for n = 1, 2,... is $\{A_t * B_t\}$.

Proof. It suffices to apply Prop. 2, to the $\sum_{n=0}^{\infty} c'_n$ and $\{C'_t\}$ defined in Prop. 3 above; since conditions (3) and (4) of the said Prop. 3 are exactly (1) and (2) of Prop. 2; and furthermore $\sum_{n=0}^{\infty} c'_n$ has the associated $\{C'_t\}$ after this Prop. 3, which is the other assumption in Prop. 2, that assures that the associated of $\sum_{n=0}^{\infty} a''_n$ is $C_t = A_t * B_t$ for $t = 0, 1, 2, \ldots$

5. Identity of the sums with the Euler's or Borel's summation algorithms, and with a weaker one satisfying the formal laws of no. 1.

It is expressed by the following theorem:

Theor. 1. Let σ be a summation algorithm weaker (Def. 3, no. 1) than that of Euler E \mid Borel B

If σ satisfies the formal laws a', b', and c' or

$$d'$$
, d' ₁,

satisfying besides the e' and f' with the specific topology for the algorithm

we have

$$\sigma \sum_{n=0}^{\infty} a_n = E \sum_{n=0}^{\infty} a_n \qquad (1) \qquad \qquad \sigma \sum_{n=0}^{\infty} a_n = B \sum_{n=0}^{\infty} a_n \qquad (1')$$

for the whole summable series σ :

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma}$$

Proof. Since

$$S_{\sigma} \subset S_{E}$$
, $S_{\sigma} \subset S_{B}$,

every summable series σ has an associated

$$\{A_t\}$$
 in E , A (t) in B ,

i. e., the vectorial space

$$\mathfrak{u}_1, \qquad \qquad | \qquad \qquad \mathsf{U}_1,$$

52 formed by the products $e^{-t} A(t)$ of e^{-t} $\{2^{-t} A_t\}$ of 2^{-t} of the associated Bof the series of S_{σ} , is contained in the vectorial space $\mathbf{U}_{\mathbf{0}}$ of the products of 2^{-t} for the associateds of all the summable series B, E, which in turn, belongs to the space (Def. 1, no. 1): U; i. e., we have $U_1 \subset U_0 \subset U$ (2) | (2') $u_1 \subset u_0 \subset u$ being $U_1 \sim S_{\sigma}$ $u_1 \sim S_\sigma$ (3) (3')that is, coordinable with $S\sigma$ But, by virtue of e', if $\sum_{n=0}^{\infty} a_n \ \epsilon \ \mathbf{S}_{\sigma},$ it is also $\sum_{n=0}^{\infty} a_n x^n \in S_{\sigma} \text{ for } x \in]0, 1[,$ then, the sum

$$\sigma \sum_{n=0}^{\infty} a_n \ x^n \text{ for } x \in]0, 1[$$

is a functional defined in $S\sigma$, at least, and depending on a parameter $x\varepsilon$ 0, 1[that is, according to

$$(3) \qquad \qquad (3')$$

a functional defined, at least, in

$$u_1$$
 | U_1

and depending on the parameter $x \in]0, 1[$ or on the

$$z \in]0, 1[$$
 $z \in]0, +\infty[$

writing

$$x = \frac{z}{2-z} \qquad \qquad x = \frac{1}{z+1}$$

That is, if we define:

$$\mathcal{F}_{z.}(2^{-t}A_t) = \frac{2}{2-z} \cdot \sigma \sum_{n=0}^{\infty} a_n \left(\frac{z}{2-z}\right)^n \quad (4) \quad \mathcal{F}_{z}\left[e^{-t} A(t)\right] = \frac{1}{z+1} \cdot \sigma \sum_{n=0}^{\infty} \frac{a_n}{(z+1)^n}$$

for the associated

$$\{A_t\}$$
 in E $A(t)$ in B

of each series

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma},$$

we have a functional defined, at least, in

$$u_1, \qquad \qquad U_1,$$

and depending on the parameter

$$z \in]0, 1[$$
 | $z \in [0, +\infty[$

The functional \mathcal{F}_z is furthermore defined through

$$(4),$$
 $(4'),$

not only for

$$u_1$$
, U_1 ,

but for the vectorial space

$$\mathfrak{u}_2 \supset \mathfrak{u}_1 \qquad \qquad (5) \mid \qquad \qquad \mathfrak{U}_2 \supset \mathfrak{U}_1 \qquad \qquad (5')$$

(wider in fact, after e') formed by the products

$$\{2^{-t} A_t\}$$
 of 2^{-t} $e^{-t} A(t)$ of e^{-t}

by the associated of the series $\sum_{n=0}^{\infty} a_n$ such that

$$\sum_{n=0}^{\infty} a_n x^n \in [S_{\sigma} \text{ for } x \in]0, 1[$$

that is

$$\sum_{n=0}^{\infty} a_n \left(\frac{z}{2-z} \right)^n \epsilon \, S_{\sigma} \text{ for } z \in [0, 1[$$

$$\sum_{n=0}^{\infty} \frac{a_n}{(z+1)^n} \epsilon \, S_{\sigma} \text{ for } z \in [0, +\infty[];$$

among which are obviously included, all the series $\sum_{n=0}^{\infty} a_n$ with $\overline{\lim}_{n\to\infty} |a_n|^{1/n} \le 1$, because of the permanency of the algorithm σ .

Furthermore, since σ verifies the laws a', b' and f' with the specific topology of this functional

$$E$$
 B

this functional \mathcal{F}_z is for each

$$z \in]0, 1[$$
 $z \in]0, +\infty[$

distributive and continuous, i. e., linear [1, 23] in $\mathbf{U_1}$ with the subordinate topology, according to (2'), in U_1 (2), in u_1 by the topology of uo, uo, which is in turn, after (Def. 2, no. 3) that one defined by the metric of But $\mathbf{U_1}$ being, after (2')(2)a vectorial subspace of the metrical space U, the functional \mathcal{F}_z can be extended [1, Theor. 2, 55] to U and we have a functional continuous in U and depending on the parameter $z \in]0, + \infty[$ $z \in]0, 1[$ wich we will design with the same notation (4') (4)Now, applying successively a', b', (4), (Prop. 1, no. 3), c'_1 and (4), a', (4'), b' (Prop. 1, no. 4), c'₁ and (4') we have: $z \, \mathcal{F}_z \left[\Delta \left(2^{-t} A_t \right) \right] = z \, \mathcal{F}_z \left(2^{-t-1} \, \Delta A_t \right) \, - \qquad \left| \qquad \mathcal{F}_z \left\{ \frac{d}{dt} \left[e^{-t} \, A \left(t \right) \right] \right\} \right. = \mathcal{F}_z \left[e^{-t} \, A' \left(t \right) \right] - \left[e^{-t} \, A' \left(t \right) \right] \right]$

$$z \, \mathcal{F}_{z} \left[\Delta \left(2^{-t} A_{t} \right) \right] = z \, \mathcal{F}_{z} \left(2^{-t-1} \Delta A_{t} \right) - \left[\mathcal{F}_{z} \left\{ \frac{d}{dt} \left[e^{-t} A \left(t \right) \right] \right\} \right] = \mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] - \left[\mathcal{F}_{z} \left[e^{-t} A' \left(t \right) \right] \right] - \left[\mathcal{F}_{z} \left[e^{$$

$$-\frac{z}{2} \mathcal{F}_{z} (2^{-t} A_{t}) = \frac{2-z}{2} \mathcal{F}_{z} (2^{-t} A_{t}) - a_{0} -$$

$$-\frac{z}{2} \mathcal{F}_{z} (2^{-t} A_{t}) = (1-z) \mathcal{F}_{z} (2^{-t} A_{t}) -$$

$$-[2^{-t} A_{t}]_{t=0}$$

$$= z \mathcal{F}_{z} [e^{-t} A (t)] - [e^{-t} A (t)]_{t=0}$$

which is the incrementation law for

$$\{2^{-t} A_t\} \in \mathfrak{u}_2$$

which is the derivation law for for $[e^{-t} A(t)] \in \mathbf{U_2}$

On the other hand, although the constant sequence $\xi_t=1$ for $t=0,\ 1,\ 2,\dots$ | constant function $\xi(t)=1$ for $0\leq t\leq +\infty$ does not belong to

and, therefore, neither

(2), to
$$\mathbf{u}_1$$
 (2'), to \mathbf{U}_1

it surely belongs to

$$\mathfrak{u}_2 \supset \mathfrak{u}_1; \qquad \qquad (5) \mid \qquad \qquad \mathfrak{U}_2 \supset \mathfrak{U}_1; \qquad \qquad (5')$$

since

$$\{2^t\}$$

is, according to

$$((3), no. 3)$$
 $((3'), no. 3)$

the associated

of the divergent series $\sum_{n=0}^{\infty}1=\infty$, but of radius 1, and we have, because of the permanency of σ

$$\mathcal{F}_{z}(\xi_{l}) = \frac{2}{2-z} \cdot \sigma \sum_{n=0}^{\infty} \left(\frac{z}{2-z}\right)^{n} =$$

$$= \frac{2}{2-z} \frac{1}{1-\frac{z}{2-z}} = \frac{1}{1-z} \text{ for } z \in [0, 1[]$$

$$= \frac{1}{z+1} \frac{1}{1-\frac{1}{z+1}} = \frac{1}{z} \text{ for } z \in [0, +\infty[],$$

which is the initial condition

The sequence $\{\xi_t^{(0)}\}\$ with

$$\xi_{t}(0) = 1$$
 and $\xi_{t}(0) = 0$ for $t = 1, 2, ...$

belongs evidently to \mathbf{u} , but it does to \mathbf{u}_2 ; since the sequence $\{2^{t} \xi_{t}(^{0})\}$ is, according to

((3), no. 3), the *E*-associated of the non convergent series $\sum_{n=0}^{\infty} (-1)^n$, with radius equal 1 too, and we have

$$\mathcal{F}_{s}(\xi_{t}(0)) = \frac{2}{2-z} \cdot \sigma \sum_{n=0}^{\infty} (-1)^{n} \left(\frac{z}{2-z}\right)^{n} =$$

$$= \frac{2}{2-z} \frac{1}{1+\frac{z}{2-z}} = 1$$

which is the second initial condition ((13), Theor. 2, no. 3, § 1).

Let us see now the convolution law. If

$$\{A_t\}$$
 and $\{B_t\}$ $A(t)$ and $B(t)$

are, respectively, the associated of

$$\sum_{n=0}^{\infty} a_n \in S_{\sigma} \text{ and } \sum_{n=0}^{\infty} b_n \in S_{\sigma}$$

and we write

$$c_n = a_n * b_n = \sum_{\nu=0}^n a_{\nu} b_{n-\nu} \text{ for } n = 0, 1, 2, \dots$$

by applying successively

(4) (Prop. 4, no. 4), b', d' and (4) | (4), (Prop. 3', no. 4), a', b' d'¹ and (4') we have

$$\mathcal{F}_{z} [(2^{-t} A_{t}) * (2^{-t} B_{t})] = \mathcal{F}_{z} [2^{-t} (A_{t} * B_{t})] = \mathcal{F}_{z} [2^{-t} (A_{t} * B_{t})] = \mathcal{F}_{z} [e^{-t} A_{t}] * [e^{-t} B_{t}]] * [e^{-t} B_{t}]] = \mathcal{F}_{z} [e^{-t} A_{t}] * [e^{-t} B_{t}]] * [e^{-t} B_{t}] * [e^{-t} B_{t}]] * [e^{-t} B_{t}] * [e^{-t} B_{t}]] * [e^{-t} B_{t}]] * [e^{-t} B_{t}] * [e^{-t} B_{t}]] * [e^{-t} B_{t}]] * [e^{-t} B_{t}]] * [e^{-t} B_{t}]] * [e^{-t} B_{t}] * [e^{-t} B_{t}]] * [e^{-t} B_{t}]]$$

$$\begin{array}{lll} & + c_{n}\left(\frac{z}{2-z}\right)^{n} + \ldots\right] + \\ & + \frac{2}{2-z}\frac{z}{2-z} \cdot \sigma \left[c_{0} + c_{1}\frac{z}{2-z} + \ldots + \right. \\ & + c_{n}\left(\frac{z}{2-z}\right)^{n} + \ldots\right] = \left(\frac{2}{2-z}\right)^{2} \cdot \sigma \left[c_{0} + c_{1}\frac{z}{2-z} + \right. \\ & + \ldots + c_{n}\left(\frac{z}{2-z}\right)^{n} + \ldots\right] = \left(\frac{2}{2-z}\right)^{2} \cdot \sigma \left[c_{0} + c_{1}\frac{z}{2-z} + \right. \\ & + \ldots + c_{n}\left(\frac{z}{2-z}\right)^{n} + \ldots\right] = \left(\frac{2}{2-z}\right)^{2} \cdot \sigma \left[a_{0}b_{0} + \right. \\ & + \left. \ldots + c_{n}\left(\frac{z}{2-z}\right)^{n} + \ldots\right] = \left(\frac{2}{2-z}\right)^{2} \cdot \sigma \left[a_{0}b_{0} + \right. \\ & + \left. \ldots + c_{n}\left(\frac{z}{2-z}\right)^{n} + \ldots\right] = \left(\frac{2}{2-z}\right)^{2} \cdot \sigma \left[a_{0}b_{0} + \left. \left(\frac{z}{2-z}\right)^{n} + \ldots\right] = \\ & + \left. \left(a_{1}*b_{1}\right)\frac{z}{2-z} + \ldots + \left(a_{n}*b_{n}\right)\left(\frac{z}{2-z}\right)^{n} + \ldots\right] = \\ & = \left(\frac{2}{2-z}\right)^{2} \cdot \sigma \left\{a_{0}b_{0} + \left[\left(a_{1}\frac{z}{2-z}\right)^{*}\left(b_{1}\frac{z}{2-z}\right)\right] + \\ & + \ldots + \left(\left[a_{n}\left(\frac{z}{2-z}\right)^{n}\right]^{*}\left[b_{n}\left(\frac{z}{2-z}\right)^{n}\right] + \ldots\right\} = \\ & = \frac{2}{2-z} \cdot \sigma \left[a_{0} - a_{1}\frac{z}{2-z} + \ldots + a_{n}\left(\frac{z}{2-z}\right)^{n} + \\ & + \ldots\right] \cdot \frac{2}{2-z} \cdot \sigma \left[b_{0} + b_{1}\frac{z}{2-z} + \ldots + \\ & + b_{n}\left(\frac{z}{2-z}\right)^{n} + \ldots\right] = \mathcal{F}_{z}\left[2^{-t}A_{t}\right] \cdot \mathcal{F}_{z}\left[2^{-t}B_{t}\right] \end{array}$$

$$= \frac{1}{(z+1)^2} \sigma \left[0 + a_0 b_0 + \left[\left(\frac{a_1}{z+1} \right) * \left(\frac{b_1}{z+1} \right) \right] + \cdots + \left[\frac{a_n}{(z+1)^n} * \frac{b_n}{(z+1)^n} \right] + \cdots \right] =$$

$$= \frac{1}{z+1} \cdot \sigma \left[a_0 + \frac{a_1}{z+1} + \cdots + \frac{a_n}{(z+1)^n} + \cdots \right] \cdot$$

$$\cdot \frac{1}{z+1} \cdot \sigma \left[b_0 + \frac{b_1}{z+1} + \cdots + \frac{b_n}{(z+1)^n} + \cdots \right] =$$

$$= \mathcal{F}_z \left[e^{-t} A_t (t) \right] \cdot \mathcal{F}_z \left[e^{-t} B (t) \right]$$

for every pair:

$$\{2^{-t}A_t\}$$
 ϵ \mathfrak{u}_2 and $\{2^{-t}B_t\}$ ϵ \mathfrak{u}_2

$$[e^{-t} A(t)] \epsilon \mathbf{U_2}$$
 and $[e^{-t} B(t)] \epsilon \mathbf{U_2}$

It does not seem easy to prove that these

incrementation

derivation

U

or convolution laws hold in the whole space

But this is not either necessary in order to apply the (Theor. 1 and 2, no. 3, § 1) of characterization; since after (Theor. 3, no. 3, § 1) we can take in said (Theor. 1 and 2, n.º 3, § 1), as a complete sequence in the space the

$$\varphi_{n, t} = 2^{-t} t^{(n+1)}$$
 for $t = 0, 1, 2, ...$ $\varphi_n(t) = e^{-t} t^n$ for $0 \le t \le +\infty$

and n = 0, 1, 2, ...

which not only belonge to u_2 , but also tu u_1 | whose product by e^{-at} belongs surely to $U_{2\bullet}$ if the strictly positive constant is strictly greater than 1, i. e., $a \in]0, 1[$

then, the

sequence
$$\{t^{(n+1)}\}$$

for each n = 0, 1, 2, ... is the associated

of a series $\sum_{i=0}^{\infty} \omega_i^{(n)}$ with the terms

 $\omega_{n+1}^{(n)}=(n+1)!$ and $\omega_{i}^{(n)}=0$ for $i\neq n+1$ which evidently converge

$$\sum_{i=0}^{\infty} \omega_{i}^{(n)} < \infty \text{ for } n = 0, 1, 2, ...$$

We saw before that the sequences $\{\xi_t\}$ and $\{\xi_t(0)\}$ belonged to \mathbf{u}_2 , and thereafter, also the sequence $\{\xi_t - \xi_t(0)\}$. Also the $\{\xi'_t\}$ is according to ((3), no. 3) the *E*-associated of the se-

ries $\sum_{n=0}^{\infty} (-1)^{n+1} 2n$ with convergence radius equal one.

On the other hand, being $\{2^t \varphi_n, t\}$ the associated of $\sum_{t=0}^{\infty} \omega_i^{(n)}$ and the difference

At 2^t φ_n , $t=2^t$ φ_n , $t+2^{t+1}$ Δ_t φ_n , t, for $t=0,\ 1,\ 2,\ \dots$ and each $n=0,\ 1,\ 2,\ \dots$, the associated of $\sum_{i=1}^{\infty}\omega_i^{(n)}$ (Prop. 1, n.º 4 according to a, b and c_1 , $\{2^t$ Δ_t φ_n , $t\}$ the associated of (*)

$$\frac{1}{2} \left[\sum_{i=1}^{\infty} \omega_i^{(n)} - \sum_{i=0}^{\infty} \omega_i^{(n)} \right] = -\omega_0^{(n)}/2 = 0$$
for each $n = 0, 1, ...$

function $e^{-at} t^n$

B

$$\omega_{i}^{(n)} = 0 \text{ for } 0 \le i \le n-1 \text{ (if } n > 0)$$

$$\omega_{i}^{(n)} = \frac{i!}{(i-n)!} (-a)^{i-n} \text{ for } i \ge n \ge 0.$$

which actually has a convergence radius $a \in]0, 1[$.

 $\lim_{i\to\infty} |\omega_i^{(n)}|^{1/i} = a \text{ for each } n=0, 1, 2, \dots$

Specially, for n=0, it results that the function $e^{-at} \in U_2$ and, therefore, the convolution law holds for the products (19') of (Theor. 2, no. 3, § 1).

$$\Delta \varphi_{n, t} = 2^{-t-1} (n+1) t^{(n} - 2^{-t-1} (t+1))^{(n+1)}$$
$$2^{t} \Delta \varphi_{n, t} = \frac{1}{2} \left[(n+1) t^{(n} - (t+1))^{(n+1)} \right]$$

which, up to the factor $\frac{1}{2}$, is the associated to the difference between the two series

 $\sum_{i=1}^{\infty} \omega_i^{(n)} \text{ and } \sum_{i=1}^{\infty} \omega_i^{(n)}, \text{ whose terms are all zero, excepted the n-th of the former and the } (n+1)-th \text{ of the second, both equal to } (n+1)!$

^(*) This can be so proved directly:

then

$$\{\Delta_t \varphi_{n, t}\} \in \mathfrak{u}_1 \subset \mathfrak{u}_2$$

Furthermore, according to (Prop. 2, no. 4), since $\sum_{i=0}^{\infty} \omega_i^{(n)}$ with $\omega_0^{(n)} = 0$ has the a $\{2^t \varphi_{n, t}\}$ with $\varphi_{n, 0} = 0$ for each n = 0, 1, 2, ...,the $\{2^{t+1} \varphi_n, t+1\}$ is the associated of $\sum_{i=0}^{\infty} \zeta_i^{(n)} < \infty \text{ with } \zeta_i^{(n)} = \omega_i^{(n)} + \omega_{i+1}^{(n)} \text{ for }$ i=0, 1, 2, ... and each n=0, 1, 2, ... that is, $\{2^i \varphi_{n,\ t+1}\}$ is the associated of $\frac{1}{2}\sum_{i=0}^{\infty}\zeta_{i}{}^{(n)}\ <\ \infty\ \ \text{for\ each}\ \ n=0,\ 1,\ 2,\ \dots;$ therefore, also

$$\{\varphi_h, t+1\} \in \mathfrak{u}_1 \subset \mathfrak{u}_2$$

The convolution law is then verified for all the products

(18), (19) and (20) of (Theor. 2, no. 3, § 1)

Once thus proved all the hypotheses of the characterization theorems (Theor. 1, 2, 3, no. 3, § 1) we have:

$$\mathcal{F}_{z}[2^{-t} A_{t}] = \sum_{t=0}^{\infty} 2^{-t} A_{t} z^{t} \text{ for } \{2^{-t} A_{t}\} \in \mathfrak{u}_{1}$$

$$\text{and } z \in]0, 1[$$

$$\mathcal{F}_{z}[e^{-t} A(t)] = \int_{0}^{\infty} e^{-t} A(t) e^{-tz} dt \text{ for } [e^{-t} A(t)] \in U_{1} \text{ and } z \in]0, +\infty[$$

or according to

$$\sigma \sum_{n=0}^{\infty} a_n \left(\frac{z}{2-z} \right)^n = \frac{2-z}{z} \sum_{t=0}^{\infty} 2^{-t} A_t z$$

for
$$z \in]0, 1[$$

$$\mathcal{F}_{z}[e^{-t} A(t)] = \int_{0}^{\infty} e^{-t} A(t) e^{-tz} dt \text{ for }$$

$$[e^{-t}A(t)] \in \mathbf{U}_1 \text{ and } z \in]0, +\infty[$$

$$\sigma \sum_{n=0}^{\infty} a_n \left(\frac{z}{2-z}\right)^n = \frac{2-z}{z} \sum_{t=0}^{\infty} 2^{-t} A_t z^t \qquad \sigma \sum_{n=0}^{\infty} \frac{a_n}{(z+1)^n} = (z+1) \int_0^{\infty} e^{-t} A(t) e^{-tz} dt$$
for $z \in [0, 1]$

whatever the σ -summable series $\sum_{n=0}^{\infty} a_n \in S_{\sigma}$, may be. Whence it follows the conclusion

(1), for
$$z \rightarrow 1^-$$
 | (1'), for $z \rightarrow 0^+$

by applying to the first member the law e', and to the second the classical Abel's theorem

> for powers' series for Laplace's transformation

The theorem is thus proved.

6. Algorithms of convergence of moments.

Def. 1. We call

summatory-

Stieltjes-

algorithm of convergence of moments, of p-type, to a permanent linear algorithm (Def. 1, no. 2), whose convergence factors are

$$\lambda_{t, n} = \alpha_{t} \frac{t^{(n+p)}}{\nu_{n+p}} \text{ for } t = 0, 1, 2, ...$$
 (1) $\lambda_{n}(t) = \alpha(t) \frac{t^{n+p}}{\mu_{n+p}} \text{ for } 0 \leq t < +\infty$ (1')

and n = 0, 1, 2, ..., where $p \ge 0$ is a fixed integer independent from t and n, and the a (t) a function of limited variation in $\{a_t\}$ any arbitrary sequence each finite interval [0, T],

such that

$$u_n = -\sum_{t=0}^{\infty} t^{(n)} \Delta a_t < \infty$$

$$\mu_n = -\int_0^{\infty} t^n d a (t)$$

for n = 0, 1, 2,...; or more general to the one defining the limit for a sequence $\{s_n\}$ as follows:

$$\nu \lim s_n = \sum_{t=0}^{\infty} \alpha(t) \sum_{n=0}^{t+p} s_n \frac{t^{(n+p)}}{v_{n+p}} \qquad \qquad \mu \lim s_n = \int_{0}^{\infty} \alpha(t) \sum_{n=0}^{\infty} s_n \frac{t^{n+p}}{\mu_{n+p}} dt$$

$$A_t = \sum_{n=0}^{t+p} s_n \frac{t^{(n+p)}}{v_{n+p}} \text{ for } t = p, \ p+1$$

is a polynomial whose values make a sequence

$$A_{t} = \sum_{n=0}^{t+p} s_{n} \frac{t^{(n+p)}}{v_{n+p}} \text{ for } t = p, \ p+1$$

$$A(t) = \sum_{n=0}^{\infty} s_{n} \frac{t^{n-p}}{\mu_{n+p}} \text{ for } 0 \le t < +\infty$$

is the function obtained by analytical continuation along the positive real semiaxis of the function holomorph in 0 defined by the series

which is called the associated for the sequence {sn} [17, 18] in the convergence algorithm. Remark. The proof of the permanence conditions [12, 195] requires a detailed study of the

sequence of polynomials
$$\left\{\sum_{t=0}^{n} \frac{t^{(n)}}{v_n}\right\}$$
 integral function $\sum_{n=0}^{\infty} \frac{t^n}{\mu_n}$

as a generalization of the

 $\{z^t\}$

exponential and the E_{α} of Mittag-Leffler functions.

Def. 2'. We will call regular sequence with respect to a μ -algorithm of Stieltjes moments $\{\mu_n\}$ to any increasing sequence with limit $+\infty$, and strictly greater than 1, in wich, furthermore, the generatrix a(t)remains structly positive, that is

$$\lim_{t \to \infty} r_t = + \infty, \ 1 < r_t \le r_{t+1} \text{ and } \alpha(r_t) > 0$$
for $t = 0, 1, 2, ...$ (2')

With this terminology we can briefly state the following

Prop. 1. Any linear convergence algorithm of summatory moments $\{v_n\}$ is reversible [1, 90].

Proof. In fact it is normal [9, 600]

Prop. 1'. Any algorithm subordinate to a linear μ -algorithm of integral moments through a regular sequence is reversible.

Proof. This means [1, 90] that for any convergent sequence $\{\eta_t\}$, there exists another $\{s_n\}$, convergent or not, but such that

$$\sum_{n=0}^{\infty} \lambda_n(r_t) \, s_n = \eta_t \text{ for } t = 0, 1, 2, \dots (3')$$

or after (1')

$$\sum_{n=0}^{\infty} \frac{s_n}{\mu_{n+p}} r_t^n = \frac{\eta_t}{\alpha(r_t) r_t^p} < \infty \text{ for } t = 0, 1, 2, \dots (4')$$

Now then, [1, 76], the necessary and sufficient condition for the system of infinite linear equations

$$\sum_{n=0}^{\infty} r_t^n \ z_n = \frac{\eta_t}{\alpha(r_t) r_t^{\rho}} \text{ for } t = 0, 1, 2, \dots (5')$$

to have a solution $\{z_n\}$ such that

$$\sum_{n=0}^{\infty} |z_n| = M < +\infty \tag{6'}$$

is that for any finite sequence of numbers h_0 , h_1 , ..., h_m , is be verified

$$\left| \sum_{t=0}^{m} h_t \frac{\eta_t}{a(r_t)r_{tt}^p} \right| \leq M \cdot \sup_{0 \leq n < +\infty} \left| \sum_{t=0}^{m} h_t r_t^n \right|$$
for $m = 0, 1, 2, \dots$ (7')

But being, by assumption, $r_t > 1$ for t = 0, 1, 2,... this upper bound of the second member is surely $+\infty$; therefore condition (7') is verified, whatever M may be, that is, the system (5') has a solution $\{z_n\}$ which verifies (6') for any prefixed constant M, and therefrom, the solution of of (4') follows by setting

$$s_n = z_n \ \mu_{n+p}$$
 for $n = 0, 1, 2, ...$

Prop. 2'. A convergence algorithm subordinate to a linear μ -algorithm of Stieltjes moments $\{\mu_n\}$ through a regular sequence $\{r_t\}$ (Def. 2').

Prop. 2. A convergence algorithm of p-type of summatory moments $\{v_n\}$

is perfect if the sequence

$$\{\alpha_t \ t^{(n+p)}\} \qquad \qquad \{\alpha \ (r_t) \ r_t^{n+p}\}$$

being p=0 a fixed integer independent from t and n, is very weakly complete with respect to u. (Def. 5, no. 2).

Proof. If, for a sequence $\{s_t\}$ holds:

$$\sum_{t=0}^{\infty} \lambda_{t, n} s_{t} = 0$$

$$\sum_{t=0}^{\infty} \lambda_{n} (r_{t}) s_{t} = 0$$

for n = 0, 1, 2, ..., we have evidently, according to

(1)
$$\frac{1}{\nu_{n+p}} \sum_{t=0}^{\infty} a_t \ t^{(n+p)} \ s_t = 0 \qquad \qquad \left| \begin{array}{cc} 1 \\ \mu_{n-p} \\ \end{array} \right|_{t=0}^{\infty} a \ (r_t) \ r_t^{n+p} \ s_t = 0$$

or

$$\sum_{t=0}^{\infty} a_t \ t^{(n+p)} s_t = 0$$

$$\sum_{t=0}^{\infty} a \ (r_t) \ r_t^{n+p} s_t = 0$$

for n = 0, 1, 2, ...

But, being the sequence

$$\{\alpha_t \ t^{(n+p)}\} \qquad \qquad \{\alpha \ (\gamma_t) \ \gamma_t^{n+p}\}$$

very weakly complete with respect \mathbf{u} , if $\{s_t\}$ ϵ \mathbf{u}_x it will be (Def. 5, no. 2, S 1):

$$s_t = 0$$
 for $t = 0, 1, 2, ...$

As, on the other hand, according to Prop. 1 above, the algorithm is reversible, it shall be perfect [1, 90], q. e. d.

Theor. 1. Let be (Def. 1, no. 6)
$$\nu$$

a linear convergence algorithm of p-type of moments

ν

$$\nu_n = \sum_{t=0}^n a_t t^{(n)} < \infty \qquad \qquad \mu_n = \int_0^\infty a(t) t^n dt < \infty$$

for $n = 0, 1, 2, \dots$ such that

| for any regular sequence $\{r_t\}$ (Def. 2'),

the sequence

$$\{a_t \ t^{(n+p)}\}, \qquad \{a \ (r_t) \ r_t^{n+p}\}$$

where $p \ge 0$ is an integer independent from t and n, is very weakly complete with respect to u (Def. 5, no. 2. § 1).

If λ is a permanent linear convergence algorithm, with matrix or continuous (Def. 1, no. 2), stronger (Def. 3, no. 1) than the

set or union (Def. 4, no. 1) of the subordinate algoritms (Def. 5, no. 2) to the μ through some regular sequences (Def. 2').

we have

$$\nu \lim s_n = \lambda \lim s_n \qquad (9) \quad \mu \lim s_n = \lambda \lim s_n \qquad (9')$$

for every sequence

$$\{s_n\} \in \mathbf{S}_{\lambda}$$
 $\{s_n\} \in \mathbf{S}_{\mu}$ λ -convergent μ -convergent

Proof. It is a consequence of Banach fundamental [Theorem 1, Theor, 12, 95] by virtue of

In fact, if λ comes from a matrix, we have assumption,

$$\nu \subset \lambda$$
 (10) | according to (Defs. 3 and 4, no. 1) μ (r_t) $\subset \lambda$ for every regular sequence { r_t }; (10')

and, if λ is continuous, for every sequence $\{r'_t\}$ ϵ $[0, +\infty]$ with limit $+\infty$, we have, according to

 $\nu \subset \lambda \ (r'_t)$

$$\mu (r_t) \subset \lambda (r'_t)$$

But

$$u$$
 | $\mu (r_t)$

being perfect, according to

for every sequence

$$\{s_n\} \in \mathbf{s}_{\nu}$$
 | $\{s_n\} \in \mathbf{s}_{\mu} \subset \mathbf{s}_{\mu(\tau_t)}$ (Cor. 1, Prop. 43, no. 2)

we have, according to Banach theorem:

$$\nu \lim s_n = \lambda \lim s_n$$
 | $\mu(r_t) \lim s_n = \lambda \lim s_n$

if λ comes from a matrix; or

$$\nu \lim s_n = \lambda (r'_t) \lim s_n$$
 | $\mu (r_t) \lim s_n = \lambda (r'_t) \lim s_n$

if λ is continuous and $\{r'_t\}$ any arbitrary sequence of positive numbers with limit $+\infty$.

Therefore (Prop. 4, no. 2) it is surely

$$\nu \lim s_n = \lambda \lim s_n$$

$$\mu(r_t) \lim s_n = \lambda \lim s_n$$
whence (Prop. 4, no. 2).
$$\mu \lim s_n = \lambda \lim s_n$$

Cor. The conclusions of

and Theor. 1 remain valid if the sequence

$$\{\alpha_t \ t^{(n+p)}\} \qquad \qquad \{\alpha \ (r_t) \ r_t^{n+p}\}$$

is complete with respect to \boldsymbol{u} (Def. 5, no. 2, § 1).

Proof. Since then it is surely very weakly complete with respect to u (Prop. 5, no. 2, § 1).

7. Identity of Euler and Borel limits with those of a stronger algorithm

Def. 1. The classic convergence algorithm of

which we will design in short by

$$E_0$$
 | B_0

to be distinguished from the summation algorithm (Def. 1, no. 4), is, after (Def. 1, no. 6) a convergence algorithm of the p=0 type of moments, with the generatrix

$$a_t = 2^{-t}$$
 for $t = 0, 1, 2, ...$ $a(t) = e^{-t}$ for $0 \le t < +\infty$

which has:

$$\mu_n = -\sum_{t=0}^{\infty} t^{(n)}, \ \Delta \alpha_t = \sum_{t=0}^{\infty} 2^{-t-1} \ t^{(n)} = n! \qquad \mu_n = -\int_0^{\infty} t^n \alpha \, d \alpha \, (t) = \int_0^{\infty} e^{-t} \, t^n \, dt = n!$$

for n = 0, 1, 2, ...

Def. We will call a convergence algorithm

$$E_1$$
, B_1 ,

to the algorithm of the p=1 type of moments with the same generatrix as

$$E_0$$
 | B

that is (Def. 1, no. 6) to the permanent linear algorithm with convergence factors

$$\lambda_{t, n} = 2^{-t} \frac{t^{(n+1)}}{(n+1)!} \text{ for } t = 0, 1, 2, ...$$
 $\lambda_{t, n} = e^{-t} \frac{t^{n+1}}{(n+1)!} \text{ for } 0 \le t \le +\infty$

and n = 0, 1, 2, ...

or more general obtained by analytical continuation of the associated

The convenience of introducing such a convergence algorithm

$$E_1 \neq E_0 \qquad \qquad B_1 \neq B_0$$

is based on the following Prop. 1, which will be used in the proof of (Theor. 1, no. 8)

Prop. 1. The summation and convergence algorithms of Euler,
$$E=E_1$$
 | Borel, $B=B_1$

and also, specially, their subordinates (Def. 3, no. 3, and Def. 5, no. 2)

$$B(r_t) = B_1(r_t)$$

through any arbitrary sequence $\{r_t\}$ of positive numbers with limit $+\infty$

Proof. [10, 338] [12, 212]. Writing

$$s_n = \sum_{\nu=0}^n a_{\nu}$$
 for $n = 0, 1, 2, ...,$

1et

$$\{A_t\}$$
 and $\{A'_t\}$ | $A(t)$ and $A_1(t)$

be the associated to the series $\sum_{n=0}^{\infty} a_n$ and to the sequence $\{s_n\}$ in the summation and convergence algotirhuns

$$E$$
 and E_1 B and B_2

respectively (Def. 1, no. 3 and Def. 1, no. 6).

$$\Delta (2^{-t} A'_t) = 2^{-t-1} [(\Delta A'_t) - A'_t] = D[e^{-t} A_1(t)] = e^{-t} [A'_1(t) - A_1(t)] = e^{-t} A_t \text{ for } t = 0, 1, 2, ...$$

$$D[e^{-t} A_1(t)] = e^{-t} [A'_1(t) - A_1(t)] = e^{-t} A_t \text{ for } 0 \le t < +\infty$$

and furthermore

$$\left[2^{-t} A_t\right]_{t=0} = 0, \qquad \left[e^{-t} A_1(t)\right]_{t=0} = 0,$$

we have

$$\frac{1}{2} \sum_{v=0}^{t} 2^{-v} A_v = 2^{-t} A'_t \text{ for } t = 0, 1, 2, \dots (1) \int_0^t e^{-v} A(v) dv = e^{-t} A_1(t) \text{ for } 0 \le t < +\infty (1')$$

whence, for $t \to +0$, it follows:

$$\sum_{n=0}^{\infty} a_n \in S_E \xrightarrow{\sum_{n=0}^{\infty} a_n} \sum_{n=0}^{\infty} a_n \in S_{E_1} \left[\sum_{n=0}^{\infty} a_n \in S_E \xrightarrow{\sum_{n=0}^{\infty} a_n} \sum_{n=0}^{\infty} a_n \in S_{E_1} \right]$$

$$E \sum_{n=0}^{\infty} = E_1 \sum_{n=0}^{\infty} a_n$$

$$B \sum_{n=0}^{\infty} a_n = B_1 \sum_{n=0}^{\infty} a_n$$

$$B \sum_{n=0}^{\infty} a_n = B_1 \sum_{n=0}^{\infty} a_n$$

which are the two conditions (1) and (2) of (Def. 5, no. 1) for

$$E$$
 and E_1

$$B$$
 and B_1

Furthermore for every sequence $\{r_t\}$ of

$$r_t > 0$$
 for $t = 0, 1, 2, ...$

 $r_t \geq 0$ for $t=0,\ 1,\ 2,\ ...$ we have, specially, after (1')

$$\int_{0}^{r_{t}} e^{-v} A(v) dv = e^{-r_{t}} A_{1}(r_{t}) \text{ for } t = 0, 1, 2, ...$$

then, if

$$\lim_{t=\infty} r_t = +\infty$$

it results for $t \to \infty$.

$$\sum_{n=0}^{\infty} a_n \in S_{B(r_t)} \xrightarrow{} \sum_{n=0}^{\infty} a_n \in S_{B_1(r_t)}$$

$$B(r_t) \sum_{n=0}^{\infty} a_n = B_1(r_t) \sum_{n=0}^{\infty} a_n$$

which are conditions (1) and (2) of the said (Def. 5, no. 1) for subordinate algorithms

$$B(r_t)$$
 and $B_1(r_t)$

through any arbitrary sequence of positive numbers with limit $+\infty$

Lemma 1. The conclusions of (Theor. 1, no. 6) and that of

(Prop. 2', no. 6)

are verified for the convergence algorithms

$$E_0$$
 and E_1

 B_0 and B_1

and furthermore, every increasing sequence $\{r_t\}$ with limit $+\infty$ and terms strictly greater than one, i.e., with

 $\lim_{t\to\infty} r_t = +\infty \text{ and } 1 < r_t \le r_{t+1} \text{ for } t=0,1,2,...$

is regular (Def. 2', no. 6) with respect to B_0 and B_1 .

Proof. By virtue of (Cor. Theor. 1, no. 6), it suffices to remember that the sequences

$$\{2^{-t-1} \ t^{(n)}\}$$
 and $\{2^{-t-1} \ t^{(n+p)}\}$

$$\{e^{-r_t} \ r_t^n\}$$
 and $\{e^{-r_t} \ r_t^{n+p}\}$

are complete with respect to u, as we saw in

and furthermore, the second condition ((2') no. 6) is evidently satisfied for every sequence $\{r_t\}$ whose terms are greater than 1, since the generatrix $e^{-t} > 0$, for t > 0.

Prop. 2. The convergence algorithms

$$E_0$$
 and E_1

Prop. 2'. Every convergence algorithm subordinated to the convergence algorithm B_0 or B_1 through a regular sequence $\{r_l\}$ (Def. 2', no. 6 and Lem. 1)

is perfect.

Proof. According to Lemma 1, this proposition is included in

(Prop. 2', no. 6)

Theor. 1. If λ is a permanent convergence algorithm, with matrix or continuous (Def. 1, no. 2), stronger (Def. 3, no. 1) than

the linear convergence algorithm of Euler, E_0 (Def. 1)

the set or union (Def. 4, no. 1) of the subordinated algorithms (Def. 5, no. 2) to the linear convergence algorithm of Borel, B₀ (Def. 1) through regular sequences {r_i} (Def. 2', no. 6),

we have

$$E_0 \lim s_n = \lambda \lim s_n$$
 | $B_0 \lim s_n = \lambda \lim s_n$

for every sequence

$$\{s_n\} \in S_{E_0}$$
 $\{s_n\} \in S_{B_0}$ E_0 -convergent. $\{s_n\} \in S_{B_0}$

Proof. It is included in (Theor. 1, no. 6) by virtue of Lemma 1.

8. Identity of the Euler and Borel sums with those of a stronger algorithm.

It is expressed by the following:

Theor. 1. If a permanent, linear algorithm of summation μ , equivalent (Def. 5 no. 1) to a linear convergence algorithm λ , is stronger than

the summation algorithm of Euler, E (Def. 1, no. 4)

the set or union (Def. 4, no. 1) of the algorithms $B(r_t)$ (Def., 5, no. 3) subordinated to the linear summation algorithm of Borel, B, through regular sequences $\{r_t\}$, (Def. 2', no. 6 and Lem. 1)

we have

$$E\sum_{n=0}^{\infty}a_n=\mu\sum_{n=0}^{\infty}a_n$$

$$B\sum_{n=0}^{\infty}a_n=\mu\sum_{n=0}^{\infty}a_n$$

for every series

$$\sum_{n=0}^{\infty} a_n \in S_E$$

$$\sum_{n=0}^{\infty} a_n \in S_B$$
E-summable
B-summable

Proof. Being, by assumption

$$E \subset \mu$$
; $B(r_t) \subset \mu$ (Def. 4, no. 1)

and by virtue of (Prop. 1, no. 7):

$$E_1 = E B_1(r_t) = B(r_t)$$

it results from (Prop. 4, no. 1):

$$E_1 \subset \lambda$$
; $B_1(r_t) \subset \lambda$;

and, since after (Lem. 1, no. 7), the (Theor. 1, no. 6) can be applied, we have

$$E_1 \lim s_n = \lambda \lim s_n$$
 | $B_1 \lim s_n = \lambda \lim s_n$

for every sequence

$$\{s_n\} \in \mathbf{S}_{E_1} \quad | \quad \{s_n\} \in \mathbf{S}_{B_1}$$

Whence, according to (Prop. 1, no. 7) and to (Def. 5, no. 1) of equivalence, it follows

$$E\sum_{n=0}^{\infty}a_n = \mu\sum_{n=0}^{\infty}a_n$$

$$B\sum_{n=0}^{\infty}a_n = \mu\sum_{n=0}^{\infty}a_n$$

for every sequence

$$\sum_{n=0}^{\infty} a_n \in SE$$

$$\sum_{n=0}^{\infty} a_n \in S_B$$

q. e. d.

Remark. The existence of permanent linear algorithm of summation not equivalent to any permanent linear algorithm of convergence was pointed out in [12, 136-137]. Therefore, it has been necessary in Theor. I the hypothesis of the equivalence $\mu = \lambda$.

9. Analysis of the results of this paper and pending problems.

The theorems in numbers 5, 7 and 8 explain the central rol plaid by the Euler's algorithms, and even more so by the Borel's one, in the summation of numerical series [7] [13].

The one in no. 5 carries a hypothesis of continuity for the sum σ with the specific topology (Def. 2, no. 4) of the algorithm

$$E$$
, B ,

which is hard to prove and requires a special study for each algorithm.

But it has just been, thanks to this law of continuity with the specific topology

together with Abel's law (no. 1, \$ 2), that we have been able to give now an affirmative answer to the Kogbetliantz's question [8, 11]; which stated, as he did, without including the said continuity law f' (no. 1) had to be answered negatively [20, 1841], even postulating simultaneously, with Kogbetliantz, the two laws c'_1 and d' for suppression of the initial term and the product, in their most restricted forms, which requires

$$\sum_{n=1}^{\infty} a_n \in S_{\sigma} \text{ or } \sum_{n=0}^{\infty} (a_n * b_n) \in S_{\sigma}$$

and are no more verified by so fundamental algorithms as that of Borel; and the same negative answer is valid by adding Abel's law, e'.

The said theorem of no. 5 completes for algorithms weaker than

that of no. 8 for algorithms stronger than

$$E$$
 the set of subordinates of B through regular sequences

It remains, therefore, the problem for those algorithms

$$E$$
 j B

nor	s tronger	than
-----	------------------	------

E the union of the subordinates to B through regular sequences

that is, for which

 $S_{\mu} \subset S_{E}, \quad S_{E} \subset S_{\mu} \quad | \quad S_{\mu} \subset S_{B}, S_{\mu} \supset \bigcup (S_{B(r_{t})}, \{r_{t}\} \text{ regular})$

Since, furthermore, this theorem of no. 8 assures only equality for the sum of the

E-summable series, B-summable series,

it remains unsolved the coincidence of the sums for the series, which being not

E-summable, B-summable,

may do it, nevertheless, with some algorithms stronger than \boldsymbol{u}

E B

This is the same problem that arised when passing from convergent series, in which coincide the sums by any permanent algorithm, to the summable ones, Ordinary convergence has been now substituted by the

summation, which is much wider, and even some other algorithms might be tried, according to the theorem no. 6.

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CORRIGENDA

THE UNIQUENESS PROBLEM IN THE THEORY OF NUMERICAL DIVERGENT SERIES AND FORMAL LAWS OF CALCULUS II

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Page 23 lines 14, 22, 24, and page 24 line 19, for (somorphism) read (somorphism).